Dynamic updates for object-oriented operating-system kernels

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Motivation

- Rising complexity of modern software
- Need for early updates and patches
- High availability demands on software
- Restarts must be avoided where possible

Solution: Dynamic Update

- Install software updates at runtime
- Update gets effective without restart
Dynamic updates must do two things:

- Update of the code section
- Transfer of the current state

Basic flow of a dynamic update is:

- Application of update at safe point
- Transfer of state information
- Invocation redirection
... in operating systems

- Operating systems must fulfill the same requirements
- ... but impose further requirements to the update system
  - No or limited runtime system
  - ...

For dynamic update support, operating system must offer:

- Updatable unit
- Safe point
- State tracking
- State transfer
- Redirection of invocations
- (Version management)
Object-oriented dynamic update

- Objects contain state information and code
- Compromise between modules and binary-rewriting

Downsides:
- Updatable code is bound to system layout
- Dependency of dynamically resolved references
- Makes many optimization impossible

```c
...-
-SimpleSignal* recv=signal;
recv->signal();
-request=sendRequest;
recvRequest->sig.signal();
...```

```
<<component>>>
MMU

PCB
pid: int
sleep()
```
Implementation with object-orientation

Dynamic update with object-orientation

- Updatable unit $\rightarrow$ Objects
- Safe point $\rightarrow$ ?
- State tracking $\rightarrow$ List of instances
- State transfer $\rightarrow$ Object replacement
- (Version management)
Interposition

- Hiding an object behind another object
  - Forward calls to the object with additional prologues / epilogues
  - Interface changes for the caller's view
  - Provide new implementations of the object's methods
- Realized via the Adapter pattern
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Hot-Swapping

- State tracking problem
- References might be in use, when object is switched
- Interpose mediator object
  - Keeps track of currently used instances
  - Blocks all new calls to the object
  - When quiescence is reached, object is replaced
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Lazy Update

- Update of all instances takes time
- Not all instances need to be updated
- Mark them for update, update on first following call
- Future calls go to the updated object
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K42 Introduction

- Research kernel developed by IBM
- Focus on Linux API and ABI support
- Strongly modularized by application of object orientation
- Event-driven with short-lived kernel threads
Dynamic update in K42

- Factories for all objects
- Factories keep track of all created instances
- Dynamic update results in:
  1. Hot-swapping of the factory
  2. (Lazy) replacement of the existing instances
K42-Style Factories

Old PCB Factory

PCB A

PCB B
K42-Style Factories

Old PCB Factory

PCB A

PCB B

New PCB Factory
K42-Style Factories

Old PCB Factory

New PCB Factory

PCB A

PCB B

PCB A'

PCB B'
K42-Style Factories

Old PCB Factory

New PCB Factory

PCB A
PCB B

PCB A'
PCB B'
PCB C
Quiescence detection

Problem: How find safe point

- K42 Threads are short lived
- Exist in generations
- Updating sets a new generation
- Wait for all threads of last generation to complete
- Then update the object
Dynamic update system is based on dynamically resolved references

- Forbids optimizations
- Severe as objects are on a fine-grained level
- Comparision of two running systems would be interesting... but is not given :(

Upsides:

- Interesting approach
- Extendable for well modularized operating systems (Linux,...)
- ...and also for application software