Event-Driven Energy Accounting for Dynamic Thermal Management

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Benefits of Dynamic Thermal Management

- Cooling laptops, desktop PCs
 - reduced power needed for cooling
 - ◆ reduced fan speed & noise
- Cooling servers, server clusters
 - ◆ no need for overprovisioning, reduced costs (space, maintenance, ...)
 - safe operation in case of cooling unit failure
- Temperature sensors

Drawbacks of Existing Approaches

- If critical temperature is reached
 - ◆ throttle the CPU (e.g. halt cycles, stop clock mechanism) or
 - ◆ stop the processor execution
- But: neglect of application-, user- or service-specific requirements due to missing online information about
 - ◆ the originator of a specific hardware activation and
 - the amount of energy consumed by that activity
- Throttling penalizes all tasks

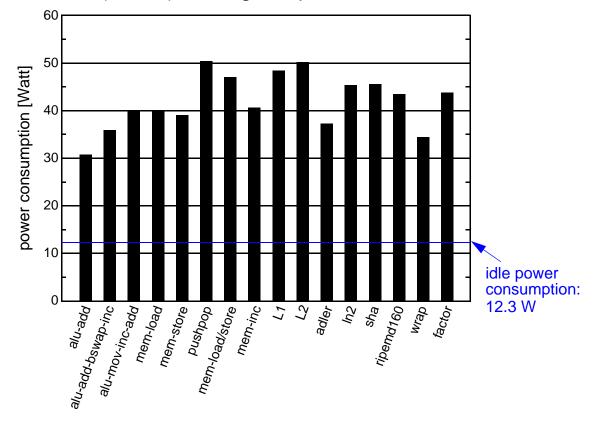
Outline

- From events to energy
 - event-monitoring counters
 - on-line estimation of energy consumption
- Energy containers
 - accounting of energy consumption
- From energy to temperature
 - temperature model
 - ◆ implementation of temperature estimation
 - ◆ task-specific temperature management

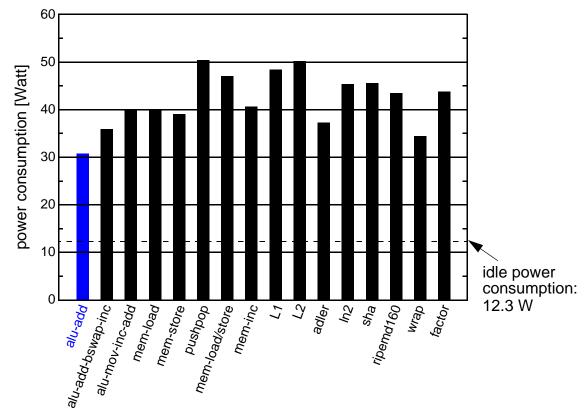
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- Counting CPU cycles
 - ◆ time as an indicator for energy consumption
 - time as an indicator for contribution to temperature level
 - throttling according to runtime
 - →but: wide variation of the active power consumption

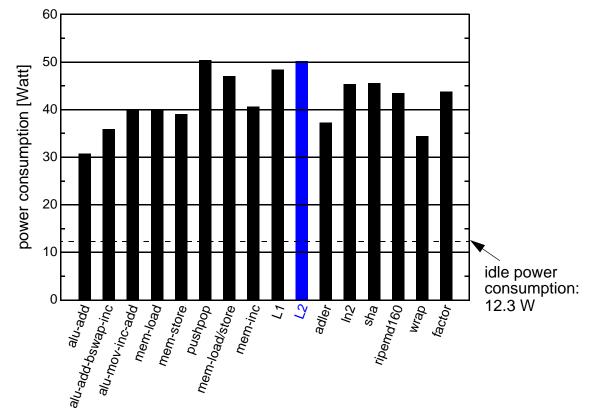
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 - →but: wide variation of the active power consumption
 - →CPU cycles are no clear indicator for energy consumption
- The case for high-resolution energy estimation

From Events to Energy: Event-Monitoring Counters

- Performance counters register energy-critical events in the complete system architecture.
 - ◆ several events can be counted simultaneously
 - low algorithmic overhead
 - ◆ low memory overhead
 - ♦ high temporal resolution
 - ◆ fast response
- Energy estimation
 - correlate a processor-internal event to an amount of energy
 - select several events and use a linear combination of these event counts
 - estimate the processor's energy consumption

From Events to Energy: Methodology

- Measure the energy consumption of training applications
- Find the events with the highest correlation to energy consumption
 - only certain combinations of events can be counted simultaneously
- Find the linear combination of these events that produce the minimum estimation error
- Avoid underestimation of energy consumption

event	weight	max. rate	power contribution
	[nJ]	(events per cycle)	@ 2GHz
time stamp counter	6.17	1.0000	12.33 W
unhalted cycles	7.12	1.0000	14.23 W
μορ queue writes	4.75	2.8430	26.99 W
retired branches	0.56	0.4738	0.53 W
mispred branches	340.46	0.0024	1.62 W
mem retired	1.73	1.1083	3.84 W
Id miss 1L retired	13.55	0.2548	6.91 W

From Events to Energy: Accuracy

- Training programs: relative error of 0–29%
- Average relative error: 5.96%
- "Real-world" applications: relative error < 6%</p>

Application	Estimation error
Mozilla 1.0.0	-0.56%
Linux 2.5.64 kernel-build	4.16%
jvm98 1.03	2.20%
caffeine 2.5	6.09%
perl	4.95%
MiBench	-1.73%

- Few cases with underestimation
 - ◆ MiBench: floating point operations
 - ◆no events for MMX, SSE & floating point instructions

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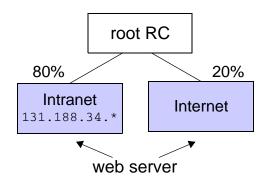
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Drawbacks of Process Accounting

- Accounting to the execution domain (processes or threads)
- Example: web server serving requests from different client classes
 - ◆ e.g. internet/intranet, different service contracts
 - cannot be distinguished on the process level
 - →accounting to different tasks/activities/clients not possible
 - "resource principal" can change dynamically

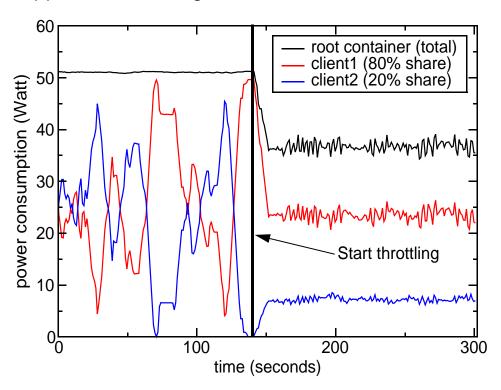
Energy Containers

- Resource Containers [OSDI '99] → Energy Containers
 - separation of protection domain and "resource principal"
- Container Hierarchy
 - ◆ root Container (whole system)
 - processes are attached to containers
- Energy shares
 - amount of energy available (depending on energy limit
 - periodically refreshed
 - if a container runs out of energy, its processes are stopped
- Packet filter attaches server to the appropriate energy container
 - depending on source address of the packet
 - energy is automatically accounted to the activity responsible for it



Energy Containers

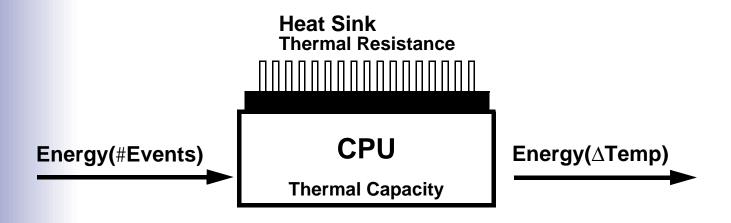
Example: server application working for two clients with different shares



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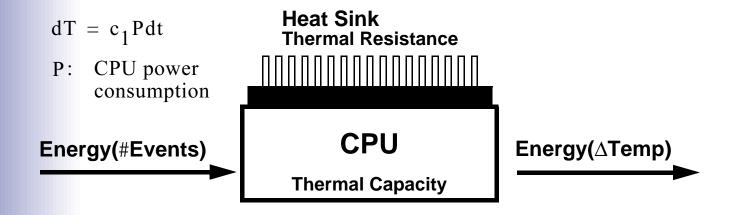
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CPU and heat sink treated as a black box with energy in- and output

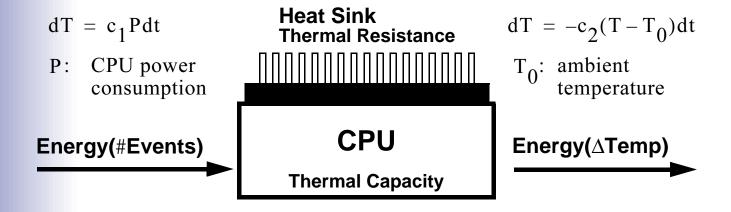


- energy input: electrical energy being consumed
- energy output: heat radiation and convection

Energy input: energy consumed by the processor



Energy output: primarily due to convection



Altogether:

$$dT = [c_1P - c_2(T - T_0)]dt$$

- ◆ energy estimator → power consumption P
- ◆ time stamp counter → time interval dt
- lacktriangle the constants c_1 , c_2 and T_0 have to be determined

Altogether:

$$dT = [c_1P - c_2(T - T_0)]dt$$

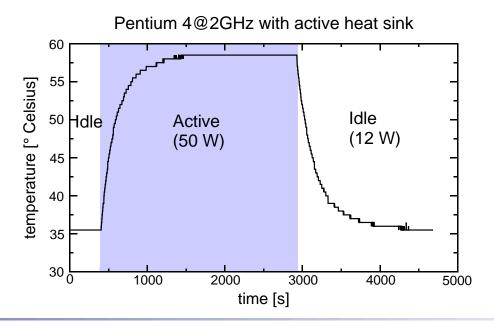
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- lacktriangle the constants c_1 , c_2 and T_0 have to be determined
- Solving this differential equation yields

$$T(t) = \frac{-c_0}{c_2} \cdot e^{-c_2 t} + \underbrace{\frac{c_1}{c_2} \cdot P + T_0}_{\text{odynamic part}}$$

$$\text{dynamic part} \quad \text{static part}$$

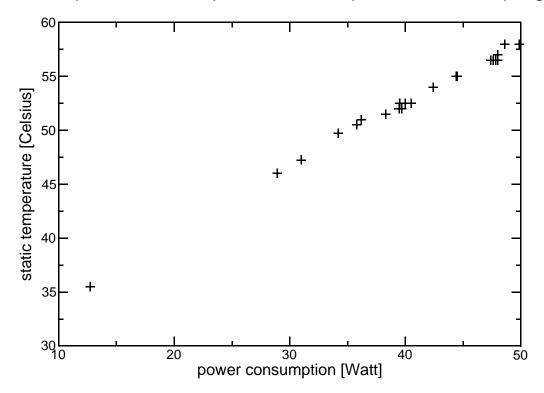
Thermal Model: Dynamic Part

- Measurements of the processor temperature
 - on a sudden constant power consumption and
 - ◆ a sudden power reduction to HLT power.
 - \rightarrow fit an exponential function to the data: coefficient = c_2



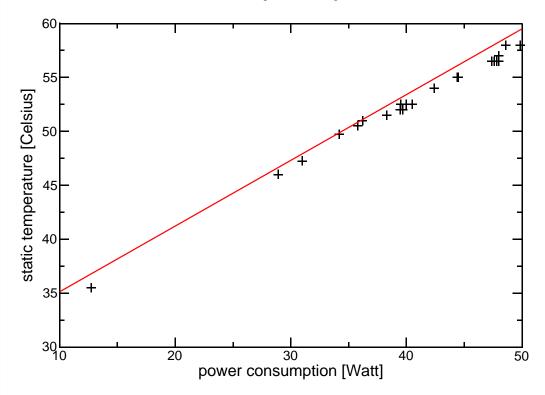
Thermal Model: Static Part

Static temperatures and power consumption of the test programs



Thermal Model: Static Part

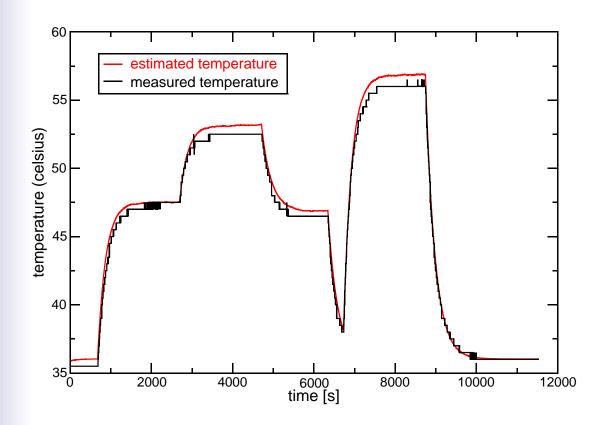
Linear function to determine c₁ and T₀



Thermal Model: Implementation

- Linux 2.5 kernel
- Periodically read the energy consumption from the root container and compute a temperature estimation
- Deviation of a few degrees celsius over 24 hours
- Re-calibration with measured temperature every 10 to 20 minutes

Thermal Model: Accuracy



Temperature Management

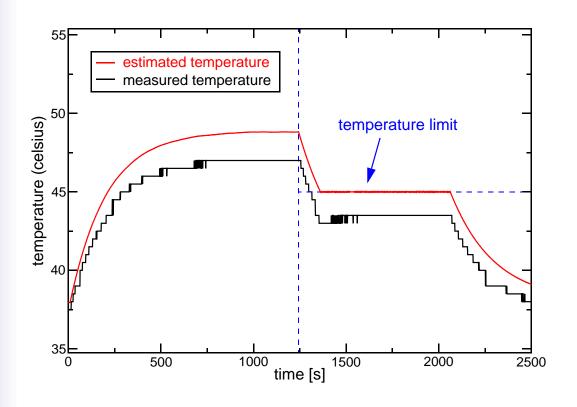
Periodically compute an energy limit for the root container (depending on the temperature limit T_{limit})

$$dT = [c_1P + c_2(T - T_0)]dt \leq T_{limit} - T$$

- Dissolve to $P \rightarrow P_{limit}$
- Energy budgets of all containers are limited according to their shares
- Tasks are automatically throttled according to their contribution to the current temperature

Temperature Management

Example: Enforcing a temperature limit of 45°



Temperature Management: Accuracy & Overhead

- Real-world applications
 - ◆ estimation within accuracy of temperature measurement (< 1° celsius)</p>
 - errors in temperature estimations always due to energy model
 - →limitations of performance counters
- Temperature limits are kept with negligible deviation
- Overhead
 - ◆ energy containers, estimation & temperature management: < 1%

Conclusion

- Event-monitoring counters enable
 - on-line energy accounting
 - ◆ task-specific temperature management
 - ◆ low overhead
- Future directions
 - examine more sophisticated energy models
 - ◆ task-specific frequency scaling to adjust the thermal load