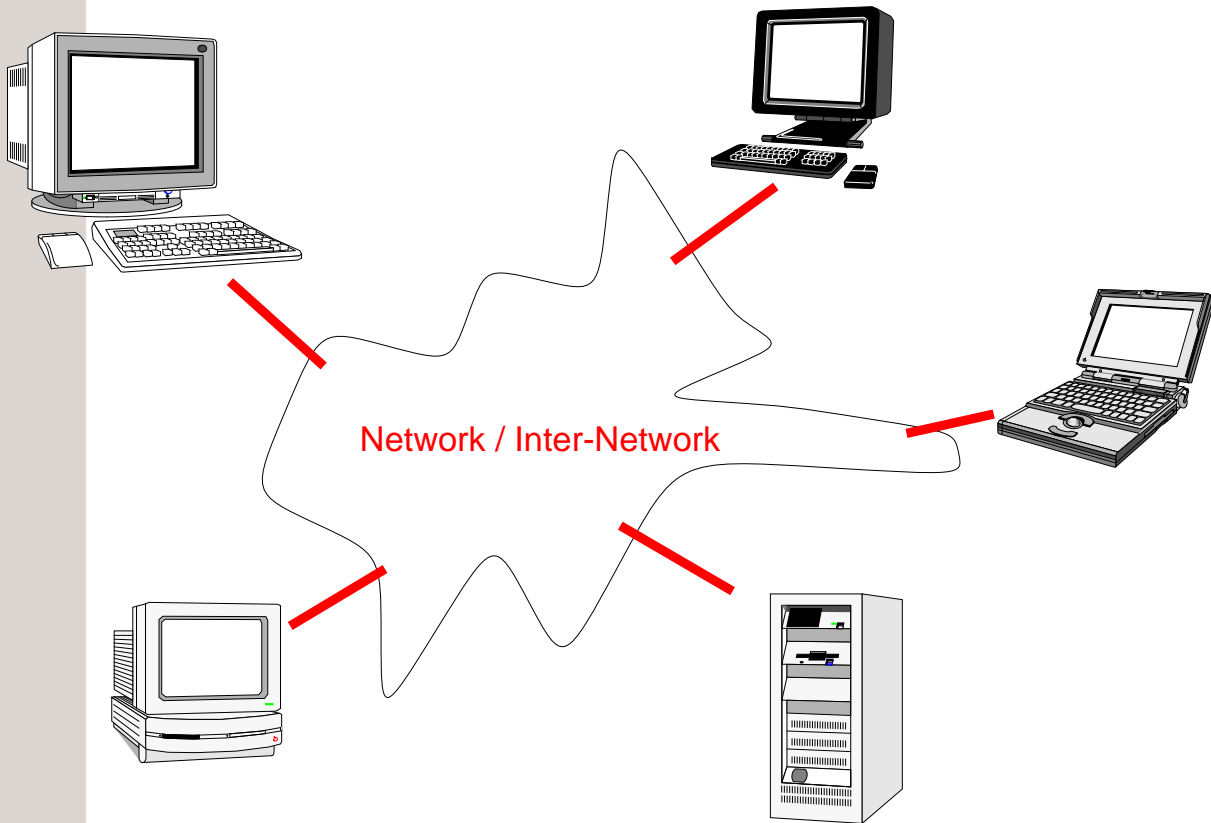


# 31 Network Programming



OODS

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Network Programming

31.218

Reproduktion jeder Art oder Verwendung dieser Unterlage bedarf der Zustimmung des Autors.

## 31.1 Host Addressing: InetAddress

- IP addresses:
  - ◆ DNS form: `www4.informatik.uni-erlangen.de`
  - ◆ "dotted quad" form: `131.188.34.42`
- `java.net.InetAddress` contains IP address
- `InetAddress` has no public constructor, create instances with
  - ◆ `getLocalHost()`
  - ◆ `getByName(String hostname)`
  - ◆ `getAllByName(String hostname)`
- convert `InetAddress` instance to different format
  - ◆ `byte[] getAddress()`: bytes of IP address
  - ◆ `String getAddress()`: "dotted quad" IP address as String
  - ◆ `String getHostName()`: host name (DNS form)

OODS

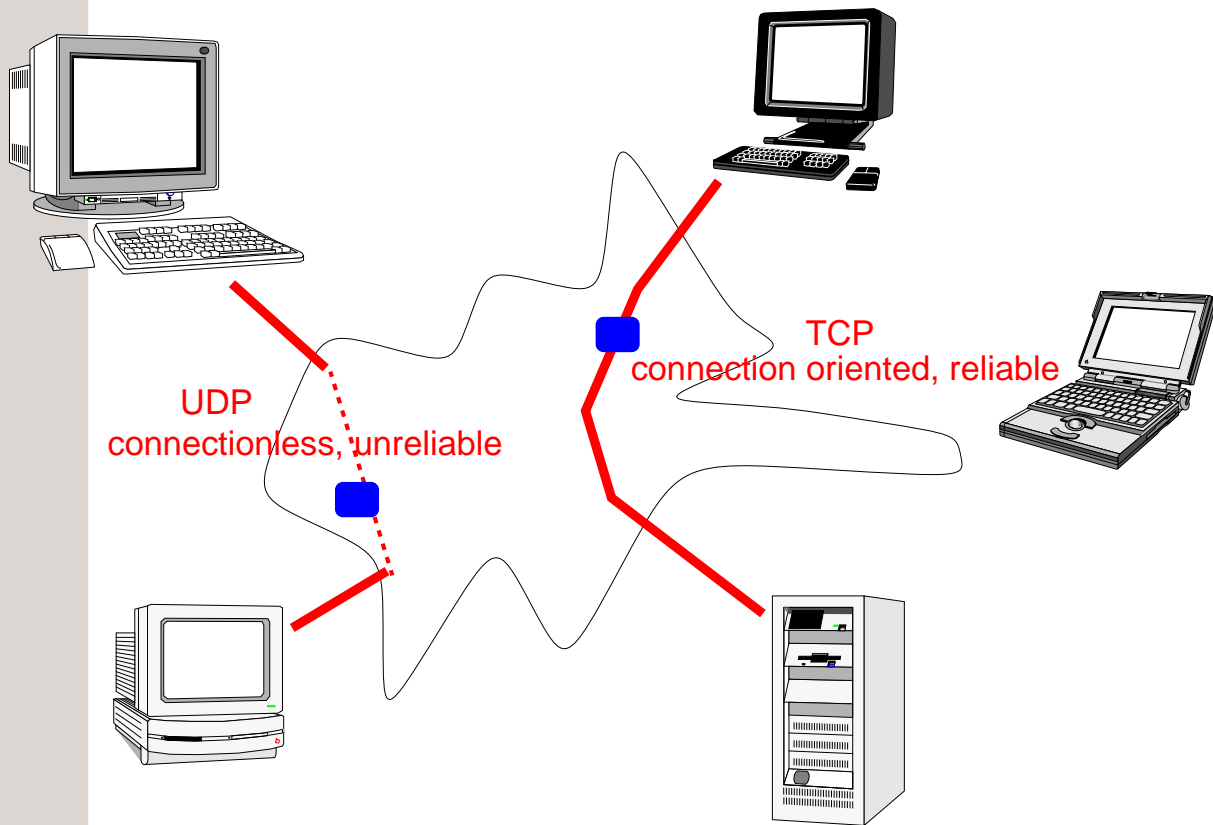
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Network Programming

31.219

Reproduktion jeder Art oder Verwendung dieser Unterlage bedarf der Zustimmung des Autors.

## 31.2 Sockets



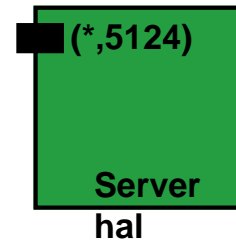
## 31.3 Connection-oriented Sockets

- `java.net.Socket`
  - ◆ TCP/IP
  - ◆ reliable
  - ◆ represents a communication endpoint at client and server
  - ◆ creating a new socket:

```
socket= new Socket("www4.informatik.uni-erlangen.de", 80);
```

- ◆ a connection endpoint is defined by *host and port* (16 bit, < 1024 privileged)
- ◆ `close` closes the socket

### 31.3.1 TCP Client/Server



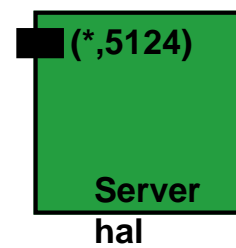
#### Server

```
ServerSocket serverSocket = new ServerSocket(5124);
```

#### Client

Empty yellow box for client code.

### 31.3.2 TCP Client/Server



#### Server

```
ServerSocket serverSocket = new ServerSocket(5124);  
Socket socket = serverSocket.accept();
```

#### Client

Empty yellow box for client code.

### 31.3.3 TCP Client/Server



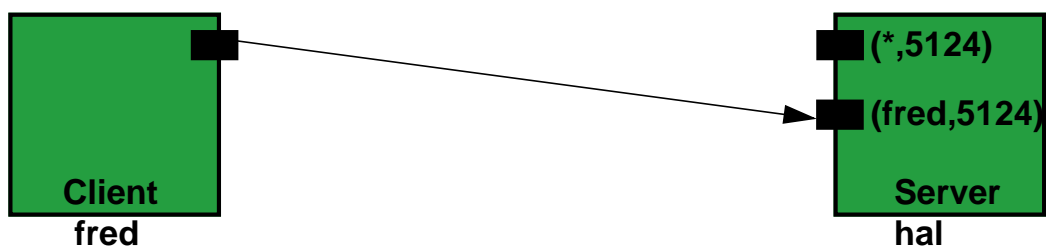
#### Server

```
ServerSocket serverSocket = new ServerSocket(5124);  
Socket socket = serverSocket.accept();
```

#### Client

```
Socket socket = new Socket("hal",5124);
```

### 31.3.4 TCP Client/Server



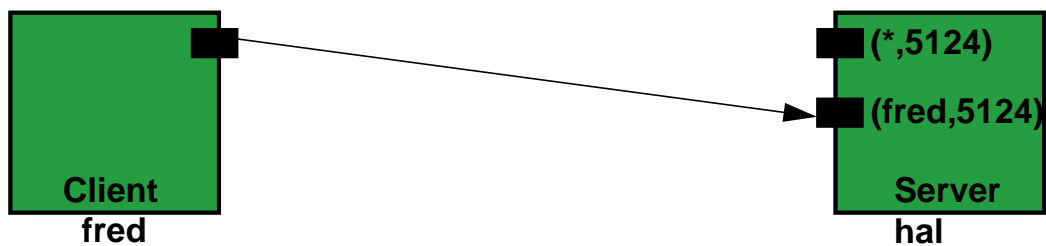
#### Server

```
ServerSocket serverSocket = new ServerSocket(5124);  
Socket socket = serverSocket.accept(); // accept returns
```

#### Client

```
Socket socket = new Socket("hal",5124);
```

## 31.3.5 TCP Client/Server



### Server

```
ServerSocket serverSocket = new ServerSocket(5124);
Socket socket = serverSocket.accept();
InputStream in = socket.getInputStream();
OutputStream out = socket.getOutputStream();
```

### Client

```
Socket socket = new Socket("hal",5124);
InputStream in = socket.getInputStream();
OutputStream out = socket.getOutputStream();
```

## 31.4 Remarks

### ■ `java.net.ServerSocket`

- ◆ is used at server side to wait for client connection requests
- ◆ `accept` waits for connection request and returns a new `Socket` object

```
ServerSocket serverSocket = new ServerSocket(10412);
Socket socket = serverSocket.accept();
```

- ◆ `close` closes the port

### ■ use streams to read/write from/to sockets

```
InputStream inStream = socket.getInputStream();
OutputStream outStream = socket.getOutputStream();
```

### ■ use these streams to create more capable streams

```
DataOutputStream out =
    new DataOutputStream(new BufferedOutputStream(outStream));
```

## 31.5 Connectionless Sockets

### ■ java.net.DatagramSocket

◆ UDP/IP

◆ unreliable: **Datagrams can get lost!**

◆ low latency

◆ Constructors:

```
DatagramSocket(int port)
    bind to local port port
```

```
DatagramSocket()
    bind to arbitrary local port
```

◆ Methods:

```
send(DatagramPacket packet)
    send packet, you must write receiver address in packet
```

```
receive(DatagramPacket packet)
    receive packet, sender address contained in packet
```

### 31.5.1 Receiver

#### ■ receive packets at specific port

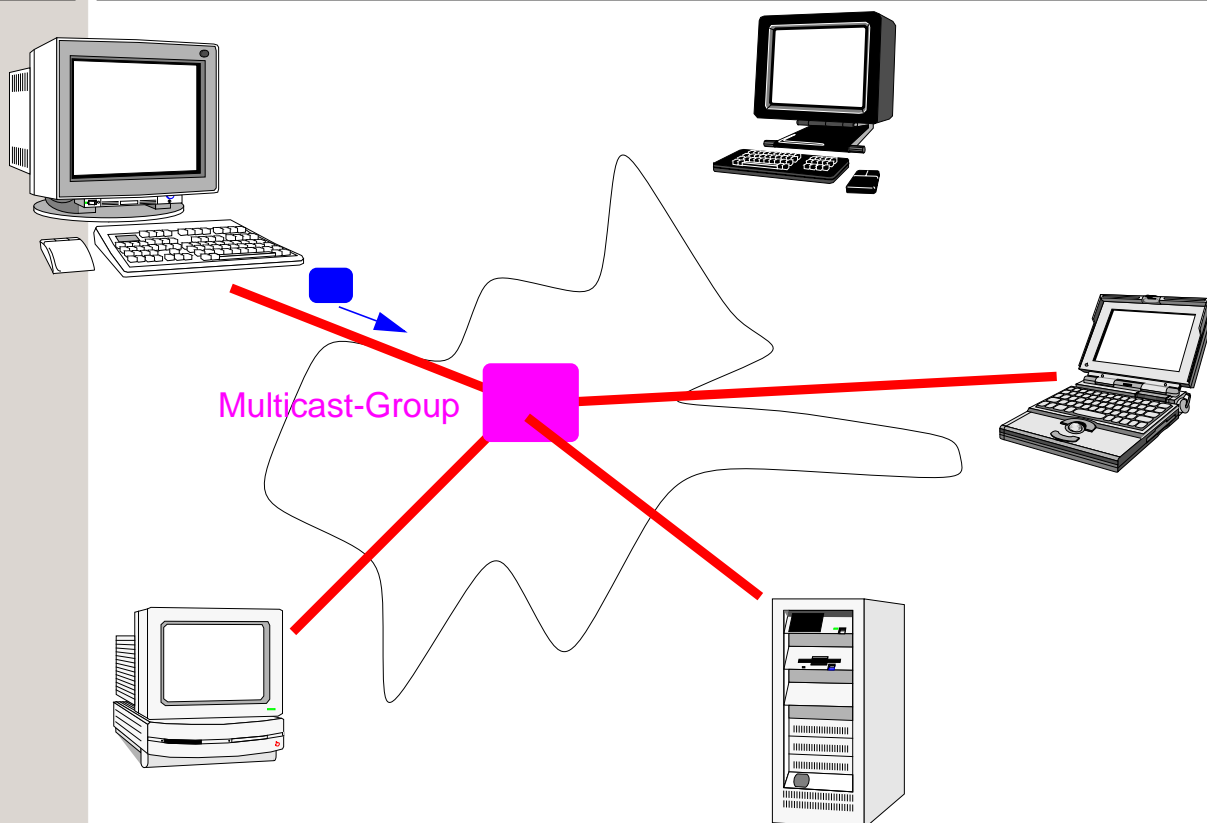
```
DatagramSocket socket = new DatagramSocket(10412);
byte[] buf = new byte[1024];
DatagramPacket packet = new DatagramPacket(buf, buf.length);
socket.receive(packet);
InetAddress from = packet.getAddress();
int bytesReceived = packet.getLength();
```

## 31.5.2 Sender

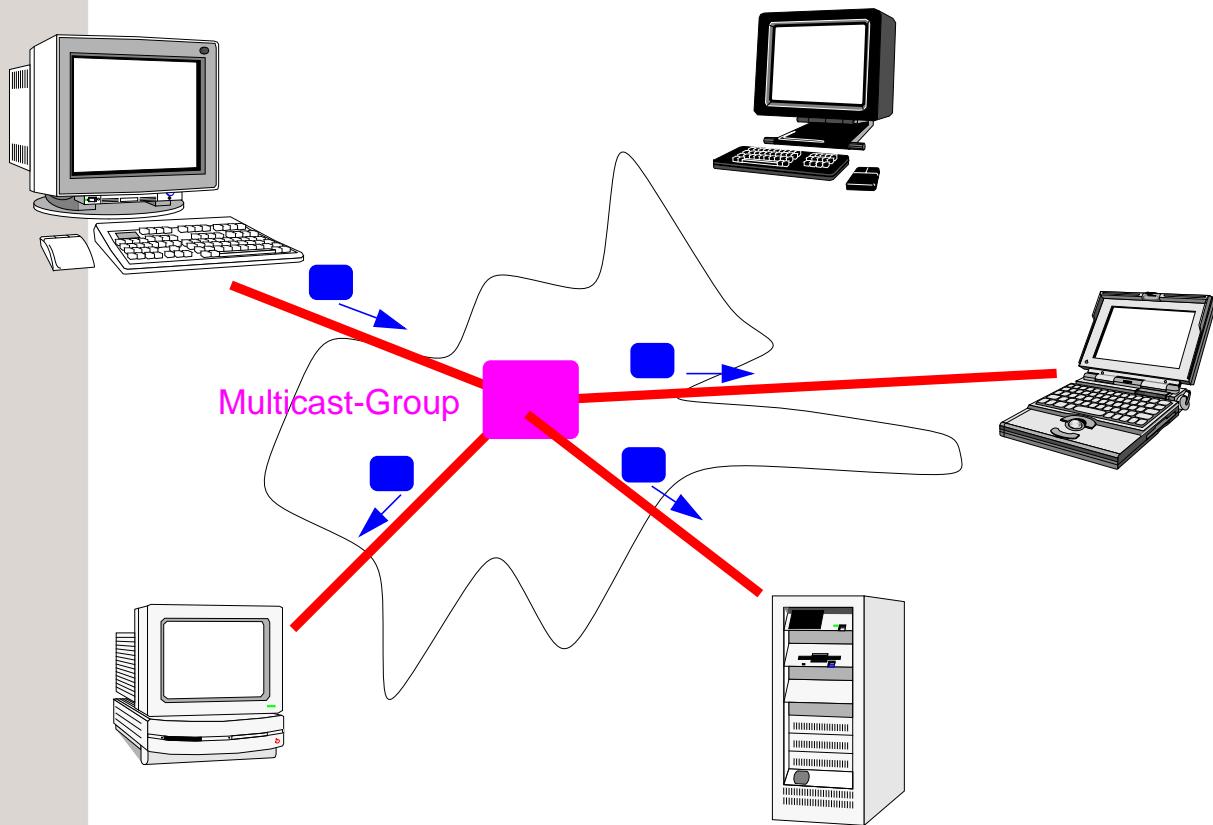
- send packets from arbitrary port

```
InetAddress addr = InetAddress.getByName("faui40");
int port = 10412;
DatagramSocket socket = new DatagramSocket();
byte[] buf = new byte[1024];
buf[0] = ...
DatagramPacket packet = new DatagramPacket(buf, buf.length,
                                           addr, port);
socket.send(packet);
```

## 31.6 Multicast-Sockets



## 31.6 Multicast-Sockets



## 31.6 Multicast-Sockets

- `java.net.MulticastSocket`
  - ◆ connection-less (subclass of `DatagramSocket`)
  - ◆ uses *class D* IP-addresses (224.0.0.1 to 239.255.255.255)
  - ◆ you can send packets after creating the socket
  - ◆ to receive packets you must join the group with `joinGroup()`
  - ◆ packet propagation is controlled by time-to-live parameter (TTL) (decremented when crossing network border)

```
InetAddress group = InetAddress.getByName("228.5.6.7");
MulticastSocket socket = new MulticastSocket(6789);
socket.setTTL((byte)2);
socket.joinGroup(group);
```

## 31.7 Applets and Sockets

---

- Security restriction:
  - ◆ applets can only open connections to the host they are loaded from
  - ◆ use `getDocumentBase().getHost()` to get the Web servers hostname

## 31.8 Summary

---

- **Socket**: one end of a two-way TCP communication (server or client)
  - ◆ contains target address/port and local port
  - ◆ use streams for reading/writing
- **SocketServer**: a server end point, creates **Socket** instances
- **DatagramSocket**: UDP communication
  - ◆ use `send()/receive()` for communication
  - ◆ target address is contained in **DatagramPacket** instance
- **MulticastSocket**: Multicast UDP communication
  - ◆ reserved range of IP-Addresses
  - ◆ use `joinGroup()` before receiving from a multicast group

## 32 Threads

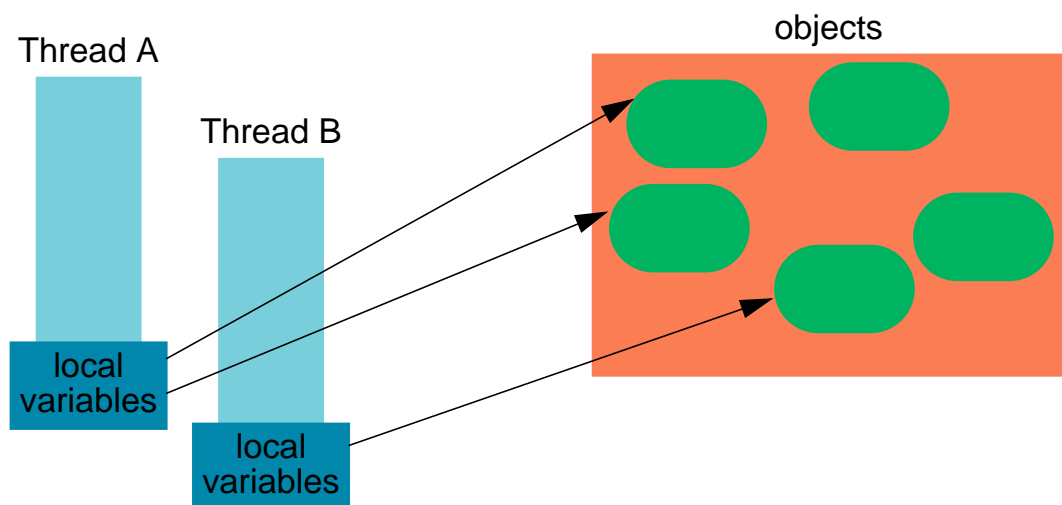
### ■ Reference

- ◆ D. Lea. *Concurrent Programming in Java - Design Principles and Patterns*. Second Edition. The Java Series. Addison-Wesley 1999.

## 32.1 Threads

### ■ What is a thread?

- ◆ unit of activity with own program counter, own registers, and own stack
- ◆ all threads use the same address space



## 32.2 Advantages / Disadvantages

### ■ Advantages

- ◆ execute parallel algorithms on a multiprocessor
- ◆ waiting for slow devices (e.g. network, user) does not block whole program

### ■ Disadvantages

- ◆ complex semantics
- ◆ difficult to debug
- ◆ John Ousterhout. *Why Threads Are A Bad Idea (for most purposes)*.  
(available from the OODS web page)

## 32.3 Thread Creation (1)

1. Subclass `java.lang.Thread` and override the `run()` method.
2. Create an instance of this class and call the `start()` method at this instance.

```
class Test extends Thread {
    public void run() {
        System.out.println("Test");
    }
}

Test test = new Test();
test.start();
```

## 32.4 Thread Creation (2)

1. Implement the interface `Runnable`, this requires implementing a `run()` method
2. Create a new `Thread` instance by passing the `Thread` constructor your `Runnable` object.
3. Call the `start()` method at the `Thread` object.

```
class Test implements Runnable {
    public void run() {
        System.out.println("Test");
    }
}
```

```
Test test = new Test();
Thread thread = new Thread(test);
thread.start();
```

## 32.5 Multithreading Problems

```
public class Test implements Runnable {
    public int a=0;
    public void run() {
        for(int i=0; i<100000; i++) {
            a = a + 1;
        }
    }
}

public static void main(String[] args) {
    Test value = new Test();
    Thread t1 = new Thread(value);
    Thread t2 = new Thread(value);
    t1.start();
    t2.start();
    try {
        t1.join();
        t2.join();
    } catch(Exception e) {
        System.out.println("Exception");
    }
    System.out.println("Expected a=200000; a="+value.a);
}
```

*What is the result of this program?*

*create two threads that use the same object*

*start the two threads*

*wait for the threads to finish*

## 32.5.1 Result

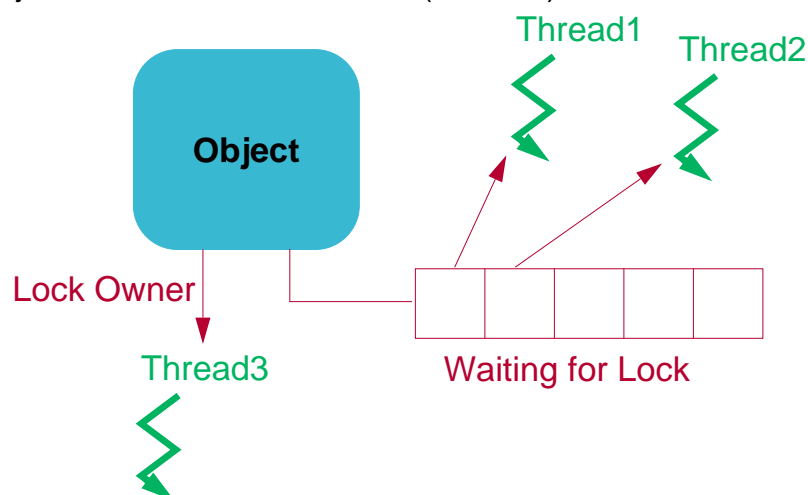
- Results of several runs: 173274, 137807, 150683
- What happens when  $a = a + 1$  is executed?

```
LOAD a into Register
ADD 1 to Register
STORE Register into a
```

- 2 possible sequences of actions when two threads are involved (initial  $a=0$ ):
  - T1-load:  $a=0$ , Reg1=0
  - T1-add:  $a=0$ , Reg1=1
  - T1-store:  $a=1$ , Reg1=1
  - T2-load:  $a=1$ , Reg2=1
  - T2-add:  $a=1$ , Reg2=2
  - T2-store:  $a=2$ , Reg2=2
  - T1-load:  $a=0$ , Reg1=0
  - T2-load:  $a=0$ , Reg2=0
  - T1-add:  $a=0$ , Reg1=1
  - T1-store:  $a=1$ , Reg1=1
  - T2-add:  $a=1$ , Reg2=1
  - T2-store:  $a=1$ , Reg2=1
- The three operations must be executed in one step (atomically)!!

## 32.5.2 Synchronization

- every object can be used as a lock (monitor)



## 32.5.3 synchronized

- use `synchronized` to acquire and release locks
- declare `method` or `block` as `synchronized`

```
class Test {  
    public synchronized void m() { ... }  
    public void n() { ...  
        synchronized(this) {  
            ...  
        }  
    }  
}
```

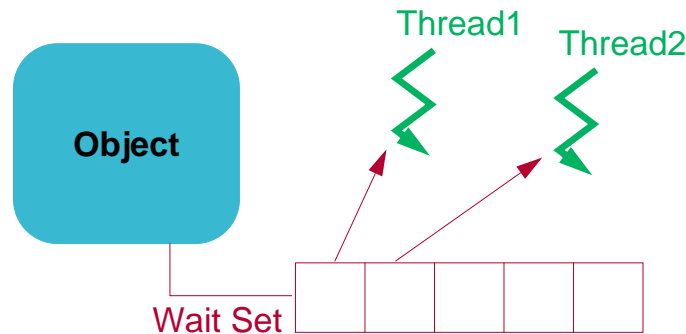
- a thread can acquire one lock multiple times (recursive lock)
- improved example: `synchronized(this) { a = a + 1; }`

## 32.5.4 Use of synchronized Methods

- no `synchronized` necessary
  - ◆ if code always runs in single-threaded context
  - ◆ for simple getter methods (see exception below)
- use `synchronized`
  - ◆ if object is written
  - ◆ if computations with the object are done (even if they *only read* the state)
  - ◆ for getter methods that read `long` or `double` types
  - ◆ for simple getter methods that must be blocked if a state update happens

## 32.6 Condition Variables

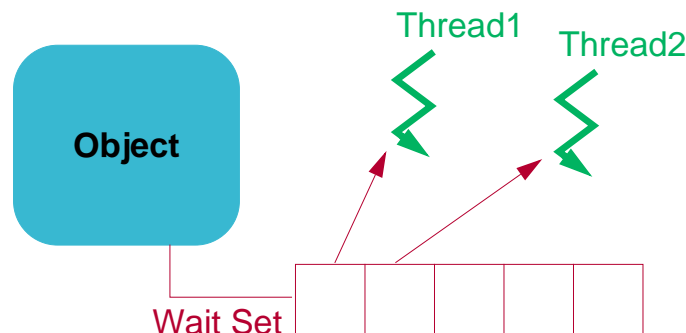
- threads must wait for a condition to become true
- two ways for waiting
  - ◆ active (polling)
  - ◆ passive (condition variables)
- every object can be used as a condition variable (every object is associated with a "wait set" = list of waiting threads)



### 32.6.1 wait and notify

- `Object` contains methods to use the object as a condition variable
  - ◆ `wait`: wait for condition (inserts thread into the wait set)  

```
while(! condition) { wait(); }
```
  - ◆ `notify`: state changed, condition could be true, inform one waiting thread (removes one thread from the wait set)
  - ◆ `notifyAll`: wake up all waiting threads (expensive!) (removes all threads from the wait set)



## 32.6.2 Condition Variables and Locks

- `wait`, `notify`, `notifyAll` can only be executed if calling thread holds object lock
- `wait` releases the lock before blocking (atomically)
  - ◆ another thread can then change the object state (making the condition true)
- acquiring the lock and unblocking happens atomically
- `wait` can be called with a timeout

## 32.6.3 Condition Variables Example

- PV system: condition is  $count > 0$

```
class Semaphore {
    private int count;
    public Semaphore(int count) { this.count = count; }
    public synchronized void P() throws InterruptedException{
        while (count <= 0) {
            wait();
        }
        count--;
    }
    public synchronized void V() {
        count++;
        notify();
    }
}
```

## 32.6.4 Condition Variables Example

- Order system: one thread accepts customer requests (`SecretaryThread`) and another thread processes them (`WorkerThread`)

```
class SecretaryThread implements Runnable {
    public void run() {
        for(;;) {
            Customer customer = customerLine.nextCustomer();
            WorkerThread worker = classify(customer);
            worker.insertCustomer(customer);
        }
    }
}

interface WorkerThread {
    public void insertCustomer(Customer c);
}
```

## 32.6.4 Condition Variables Example

- Worker

```
class SpecificWorker implements Runnable, WorkerThread {
    public void run() {
        for(;;) {
            Customer customer;
            synchronized (this) {
                while(queue.empty()) wait();
                customer = queue.next();
            }
            // do something nice with customer
            // ...
        }
    }
    public synchronized void insertCustomer(Customer c) {
        queue.insert(c);
        notify();
    }
}
```

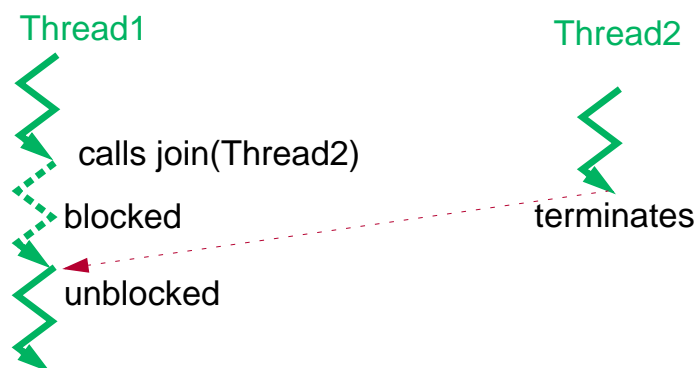
## 32.7 sleep

- thread has `sleep(long n)` method to go to sleep for n milliseconds
- thread can be suspended after return from `sleep()`
- sleep method is a static method of Thread; called like `Thread.sleep(...);`

## 32.8 join

- a thread can wait for another thread to die

```
workerThread = new Thread(worker);  
...  
workerThread.join();  
worker.result();
```



## 32.9 Daemon Threads

---

- daemon threads are used for background tasks
- should not be used for main task of program
- if all *non-daemon* threads are dead the program is finished
- How can you tell if a thread should be daemon thread?
  - ◆ You cannot state a termination condition for the thread.
- Important methods of the **Thread** class:
  - ◆ `setDaemon(boolean switch)`: turns daemon property on/off
  - ◆ `boolean isDaemon()`: tests if thread is a daemon

## 32.10 ThreadGroup

---

- Group of related threads (**ThreadGroup**):
  - ◆ Can contain threads and other thread groups.
  - ◆ Every thread can only influence threads in its own ThreadGroup
- Methods that can only used with threads of own thread group:
  - ◆ `list()`
  - ◆ `stop()`
  - ◆ `suspend()`
  - ◆ `resume()`

## 32.11 Thread States

---

New

## 32.11 Thread States

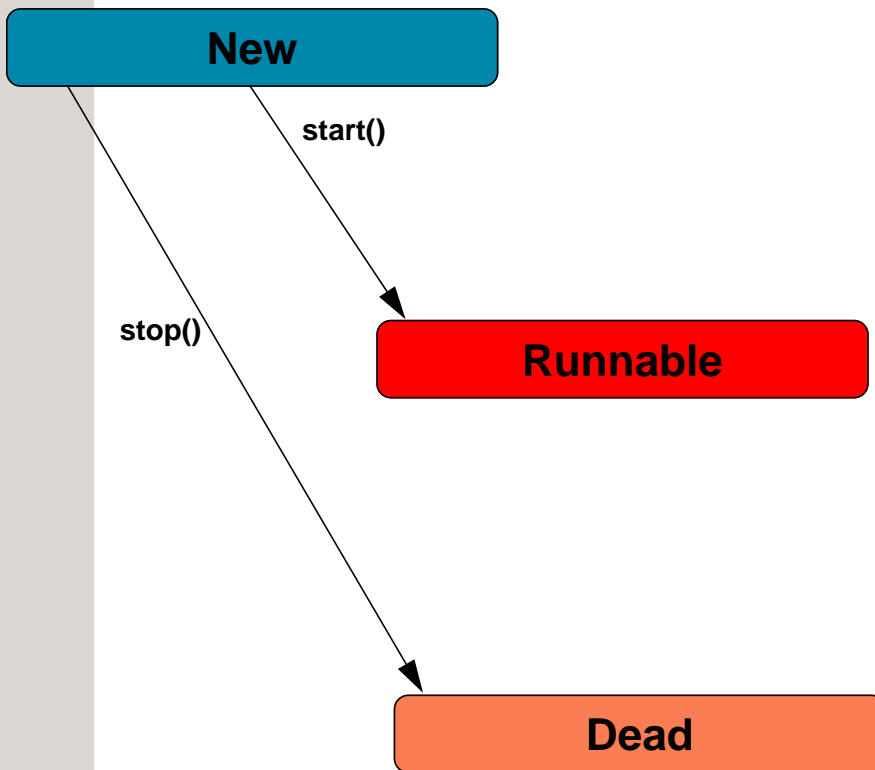
---

New

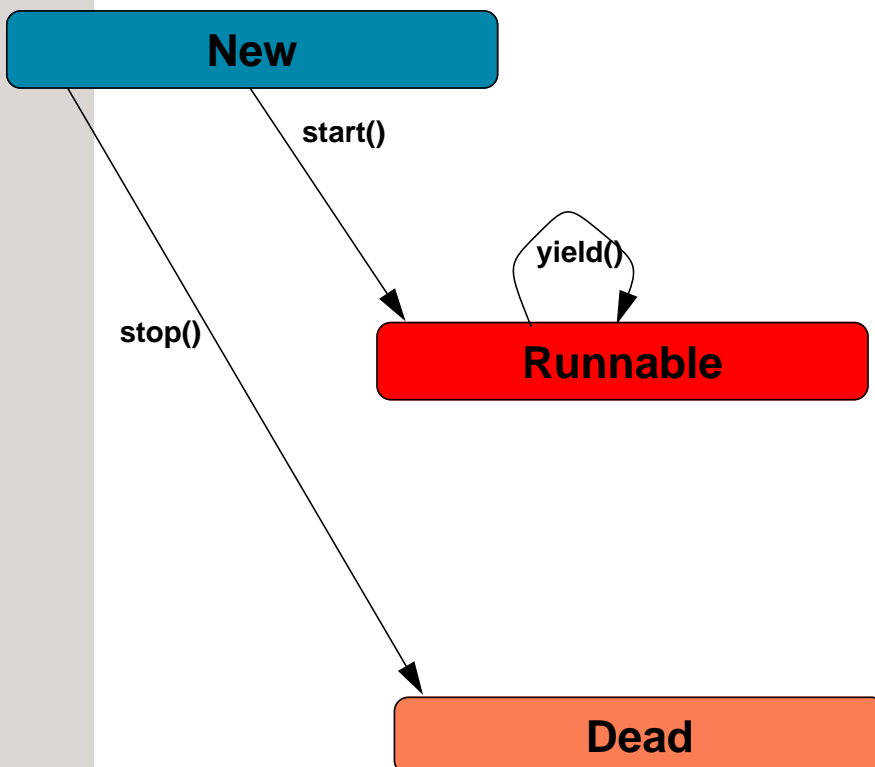
stop()

Dead

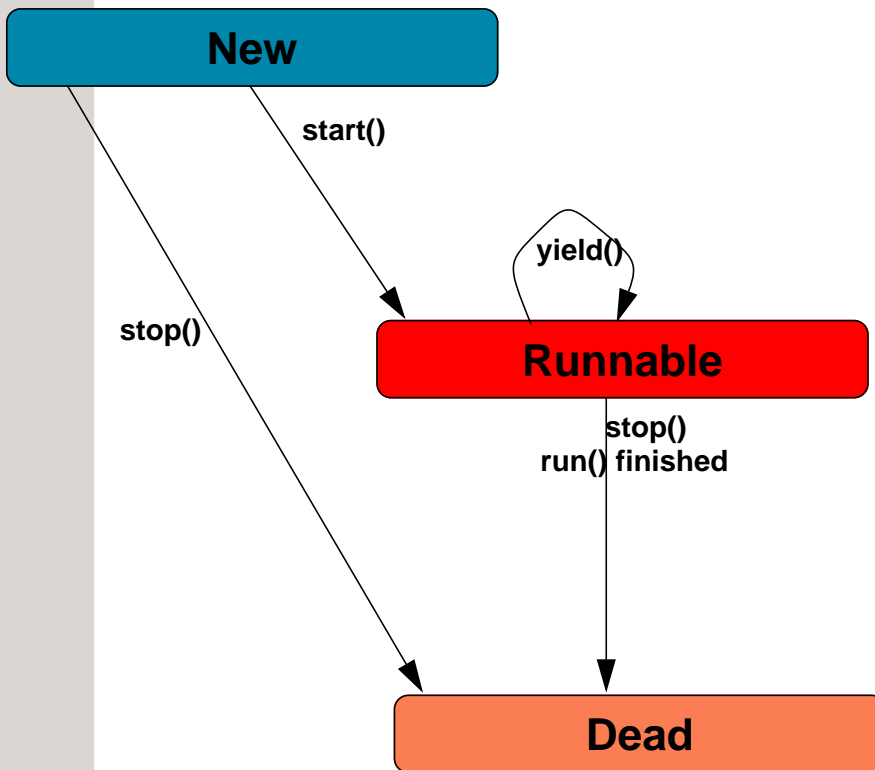
## 32.11 Thread States



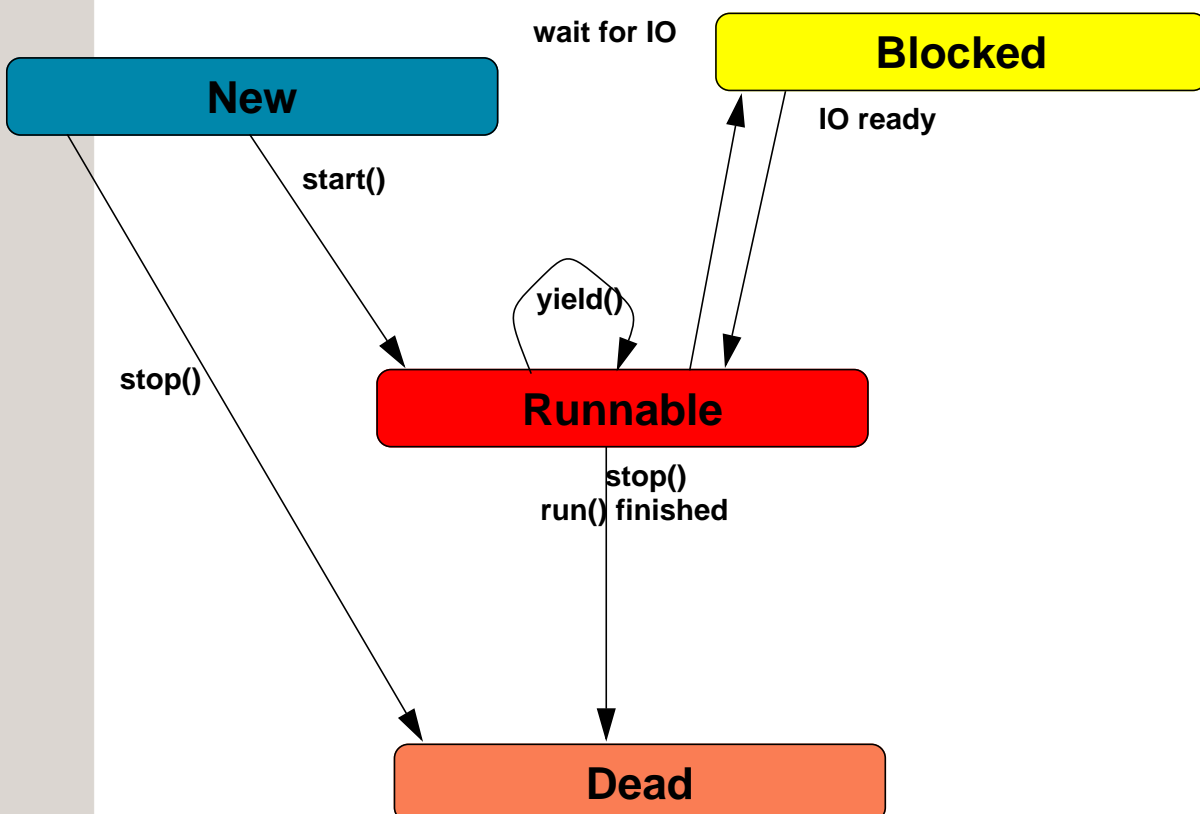
## 32.11 Thread States



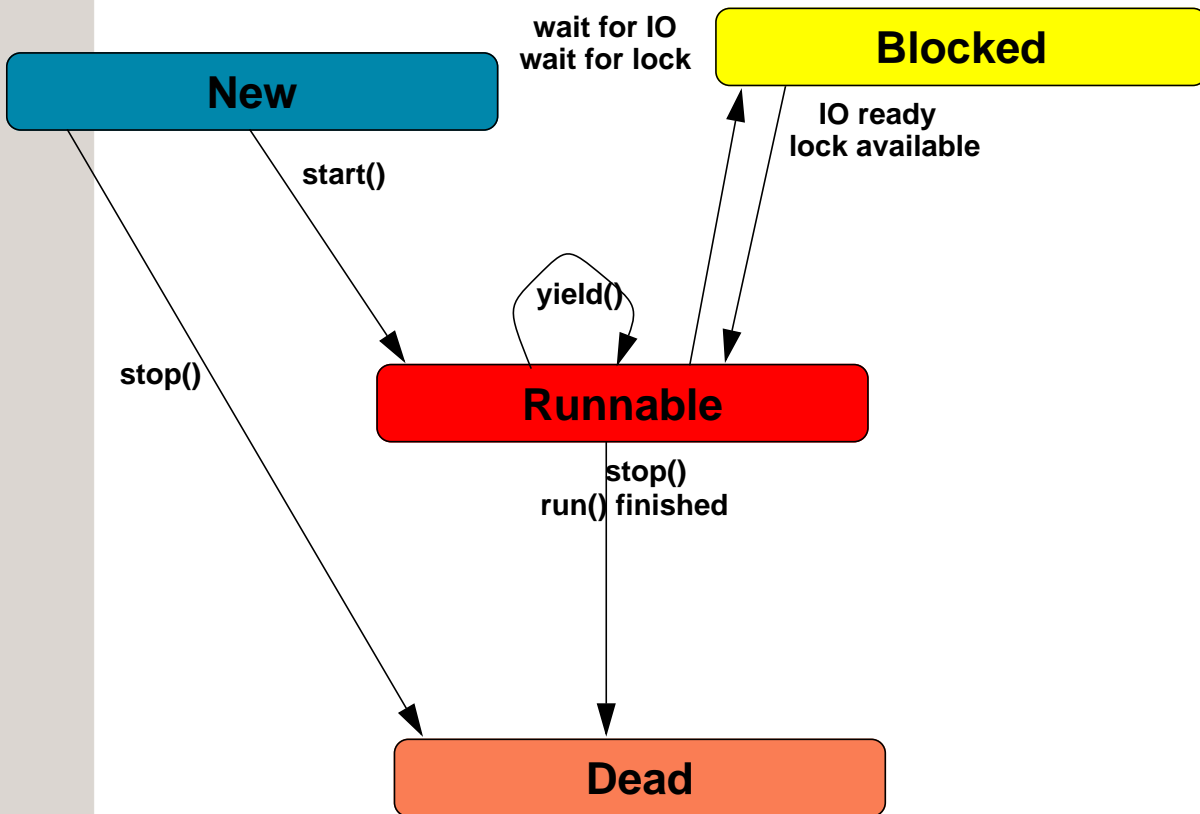
## 32.11 Thread States



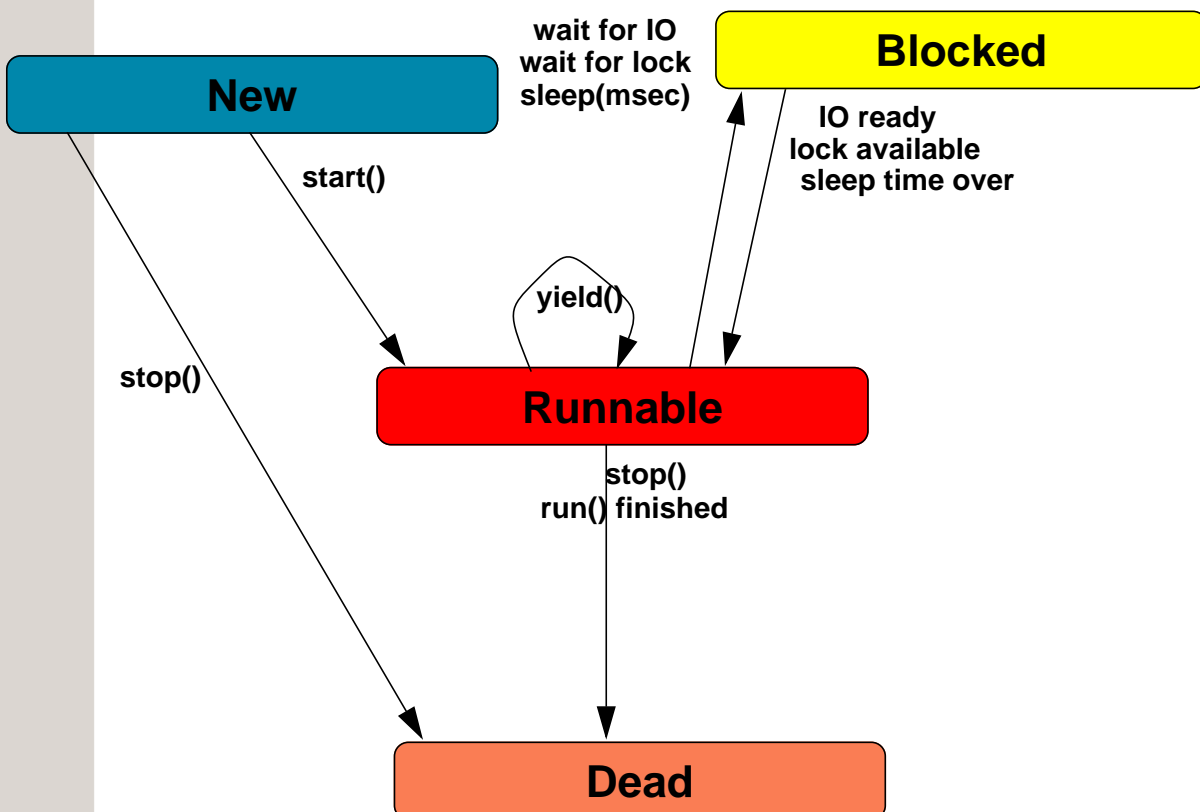
## 32.11 Thread States



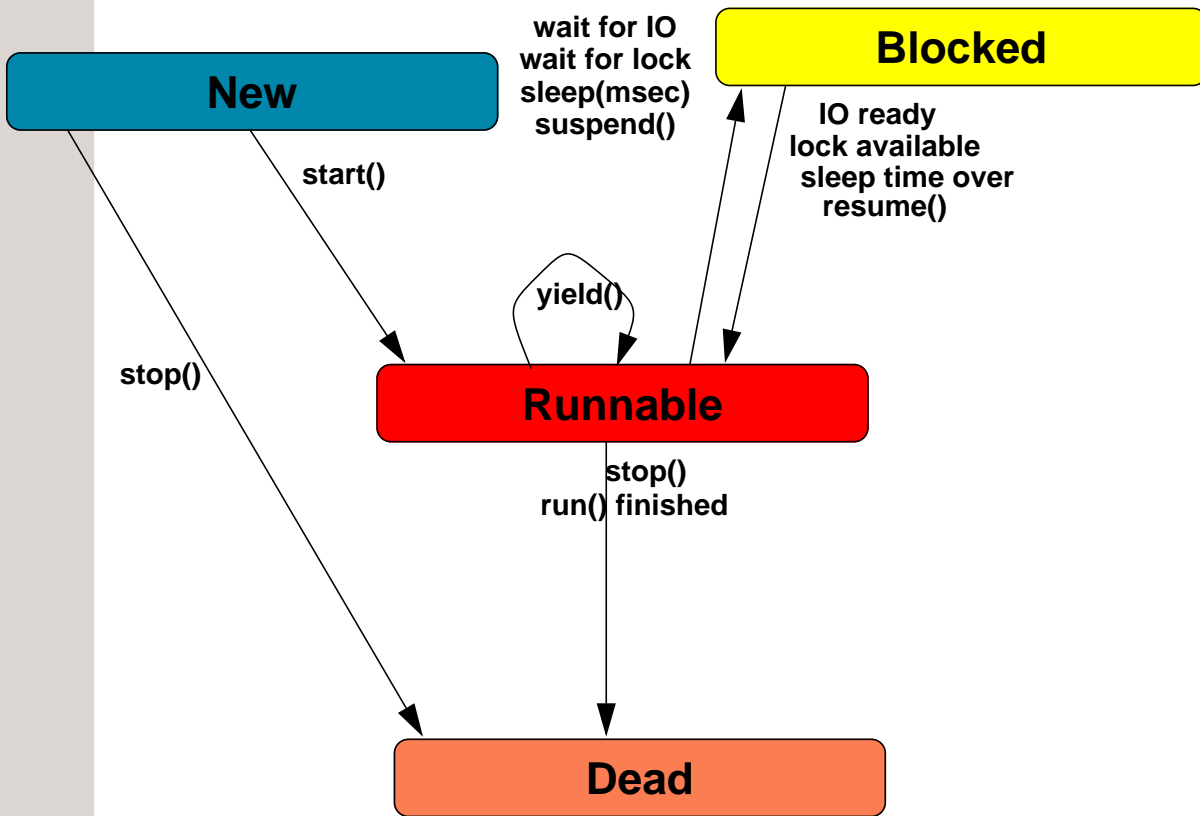
## 32.11 Thread States



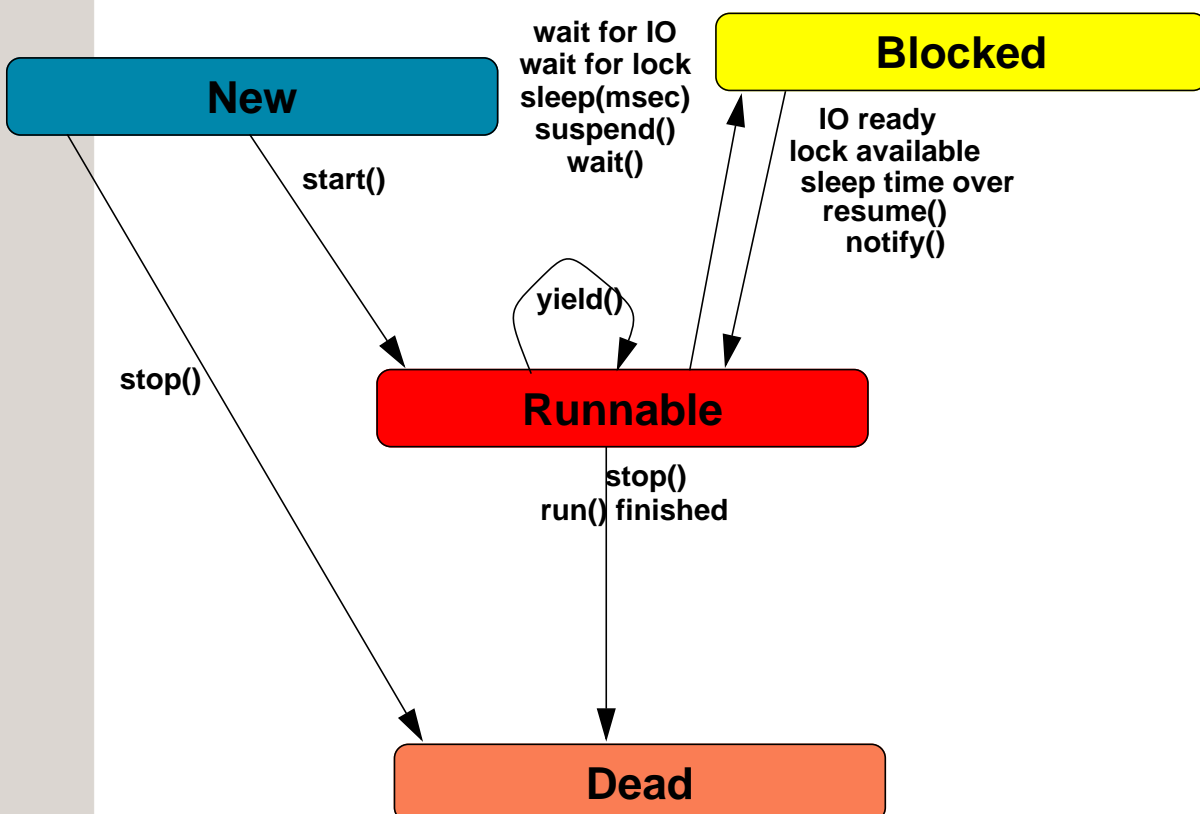
## 32.11 Thread States



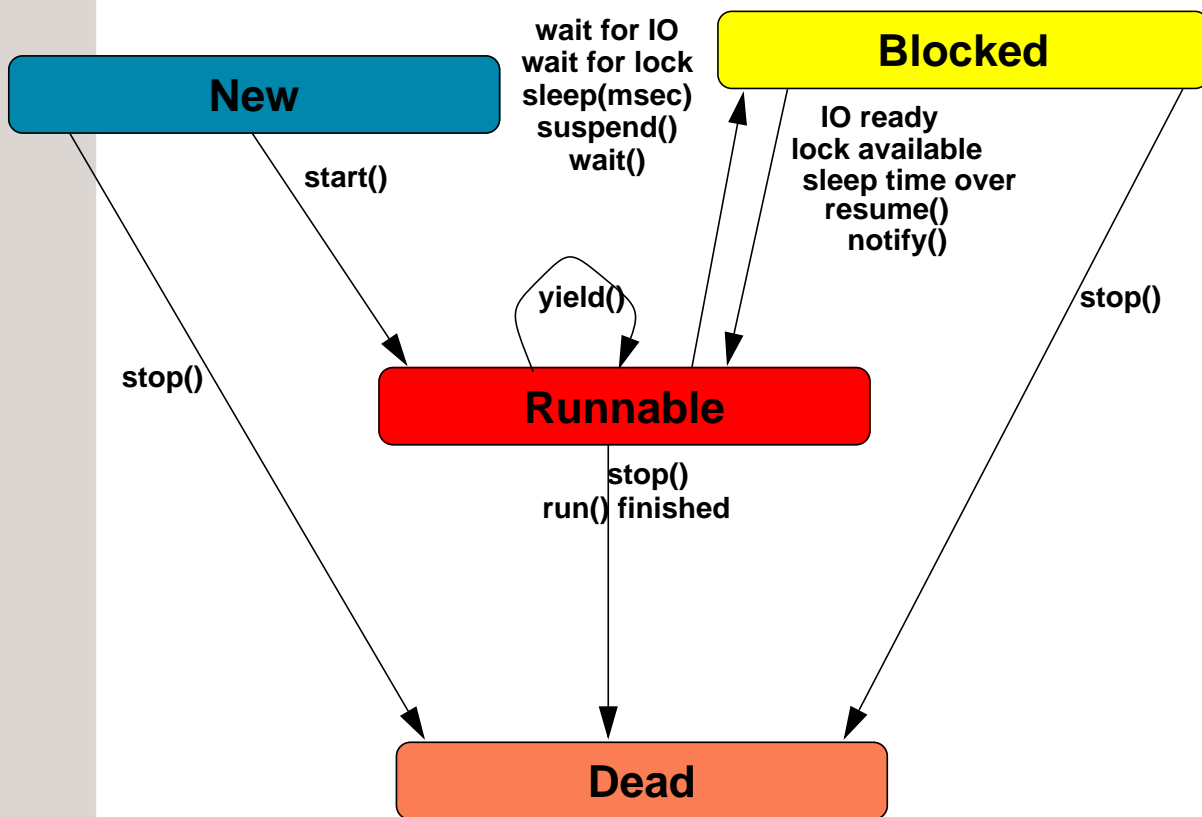
## 32.11 Thread States



## 32.11 Thread States



## 32.11 Thread States

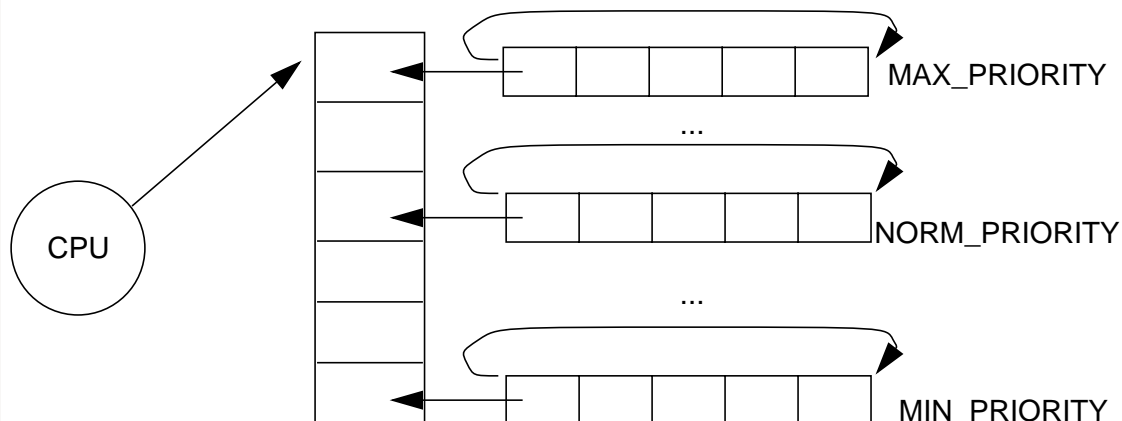


## 32.12 Scheduling

- default scheduling in jdk 1.1 on Solaris: preemptive, without *timeslicing*
  - ◆ a thread with higher priority preempts a thread with lower one
  - ◆ threads with equal priority are scheduled non-preemptive (FIFO)
  - ◆ if a thread blocks the next runnable thread with equal priority gets the CPU
- Java 1.1 and 1.2 on Solaris:
  - ◆ default or `THREADS_FLAG=green`: preemptive without timeslicing
  - ◆ `THREADS_FLAG=native`: preemptive with timeslicing
- Java 1.1 on WinNT/Win95
  - ◆ preemptive with timeslicing
- **Correctness of a program must not depend on the scheduling strategy!!!**

## 32.13 Priority based Round-Robin Scheduling

- Run queues of a timesliced, round-robin scheduler with static priorities



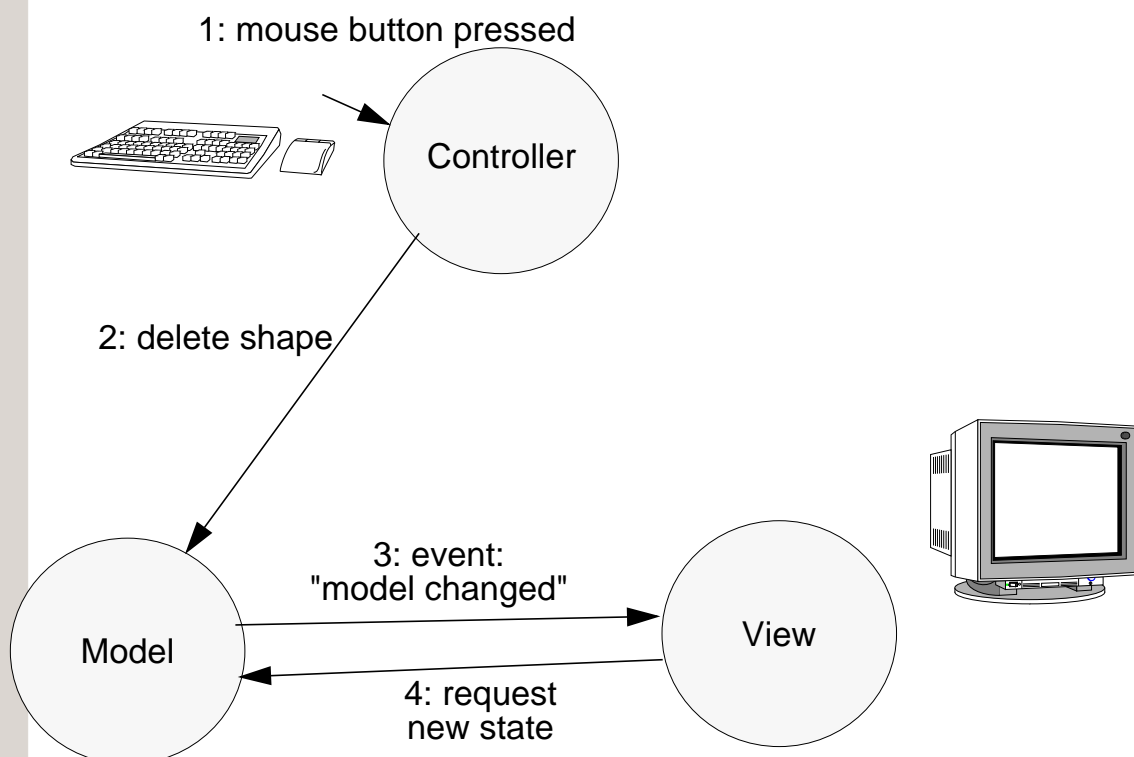
## 32.14 Deprecated Methods of Thread

- `stop()`, `suspend()`, `resume()` are deprecated in Java 1.2
- `stop()` releases all locks the thread holds - this is unsafe
- `suspend()` and `resume()` could lead to deadlock - stopped thread holds locks

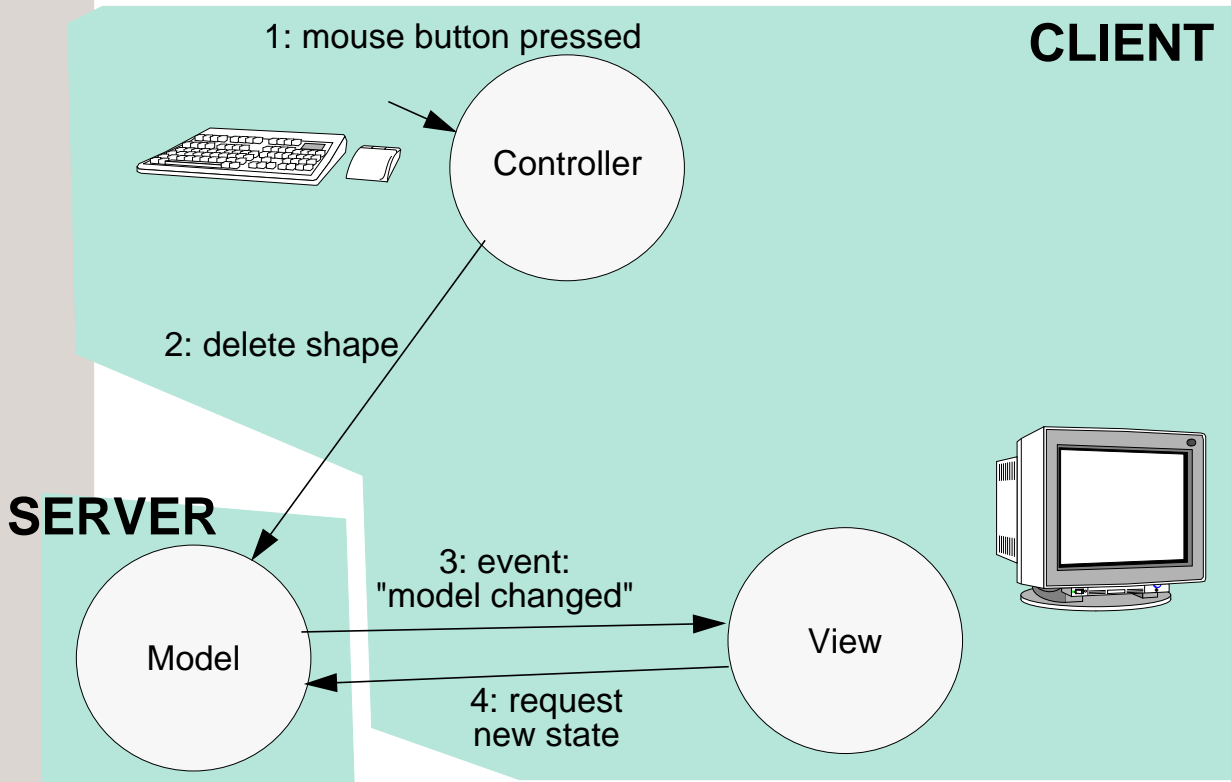
## 33 Design Patterns for the C/S Assignment

- "elements of reusable design"
- Examples
  - ◆ Model-View-Controller
  - ◆ Observer
  - ◆ Iterator
  - ◆ Command
  - ◆ Proxy

### 33.1 Model-View-Controller

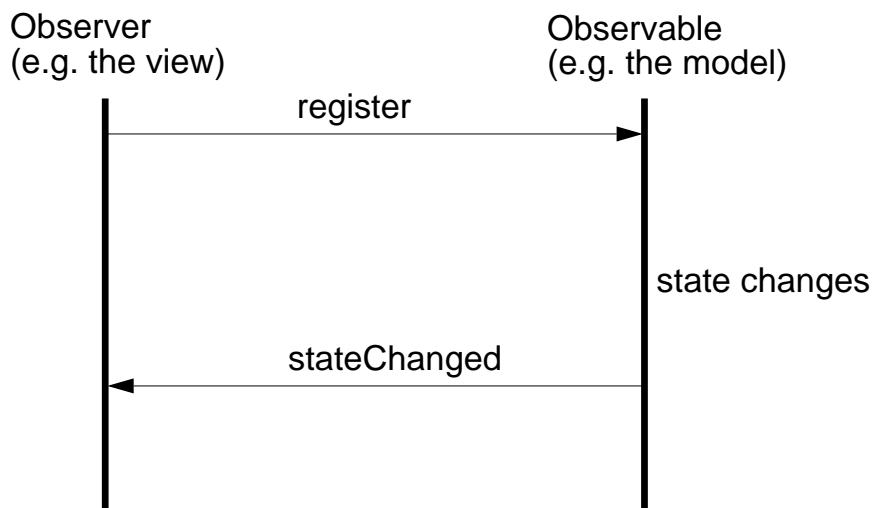


## 33.1 Model-View-Controller



## 33.2 Observer

- used in MVC by the view to observe model changes



## 33.3 Iterator

- used to "walk through" a set of objects
- iterator is responsible for maintaining the current position

```
class Iter implements java.util.Enumeration {
    int cursor;
    Shape[] shapes;

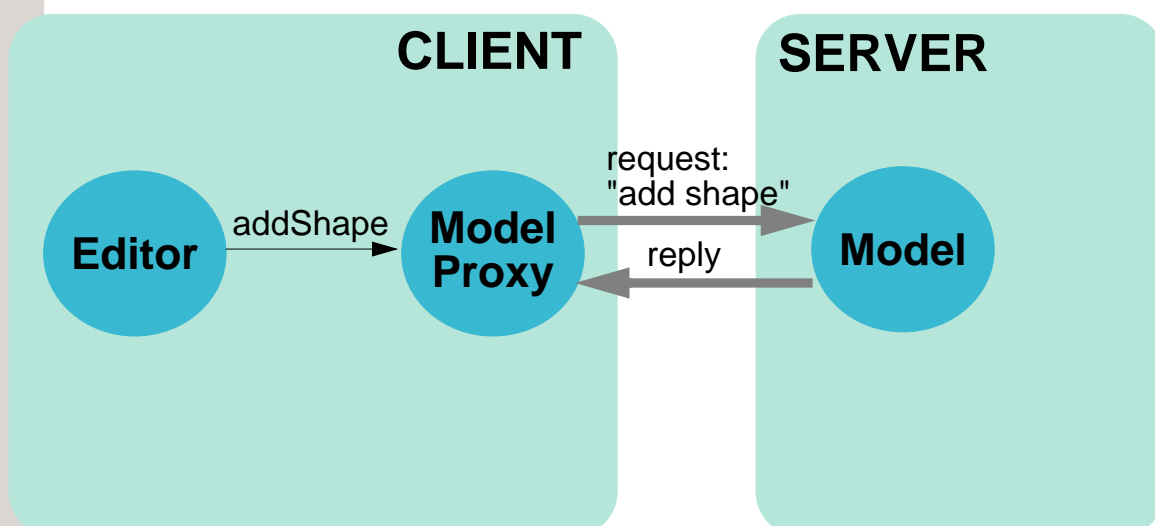
    public Iter(Shape[] shapes) {
        this.shapes = new Shape[shapes.length];
        System.arraycopy(shapes, 0, this.shapes, 0, shapes.length);
    }

    public boolean hasMoreElements() {
        while (cursor < shapes.length && shapes[cursor] == null) cursor++;
        return cursor < shapes.length; }

    public Object nextElement() { return shapes[cursor++]; }
}
```

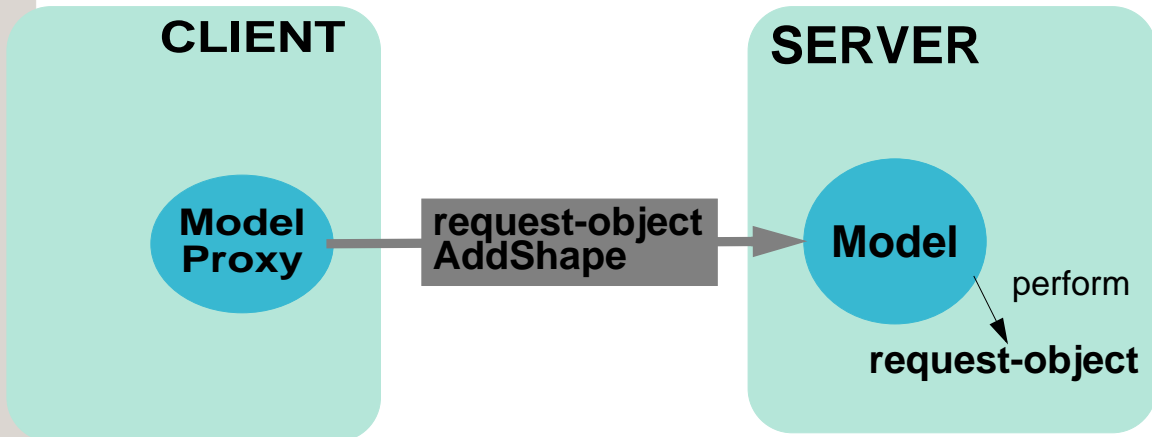
## 33.4 Proxy

- can be used to have a local representative of a remote object
- implements the same interface as the "real" object
- DrawingEditor with remote model:



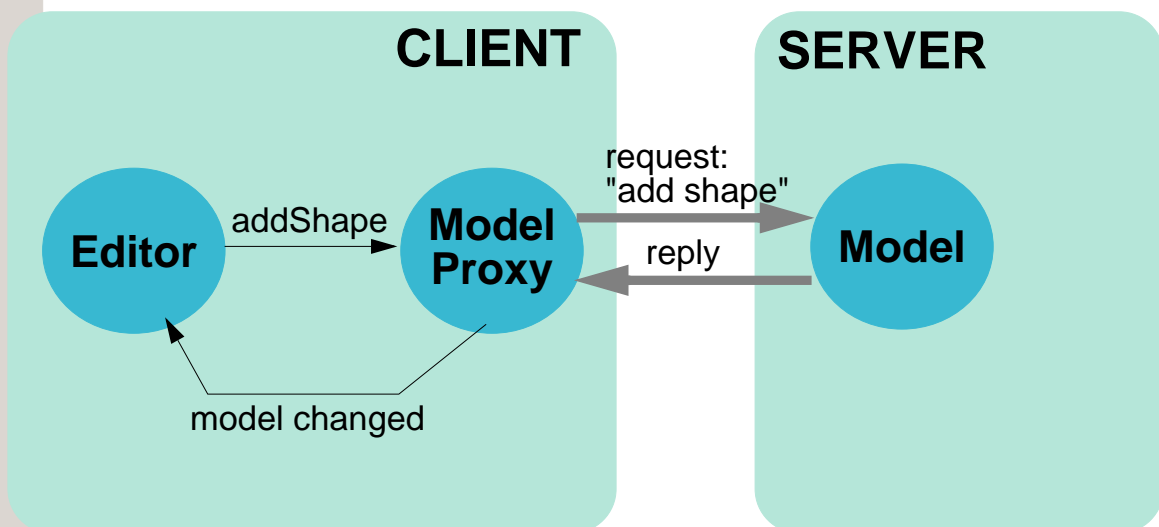
## 33.5 Command

- used to transfer a request to the server
- state contains information from client
- method perform() is called by server
- parameters contain information from server



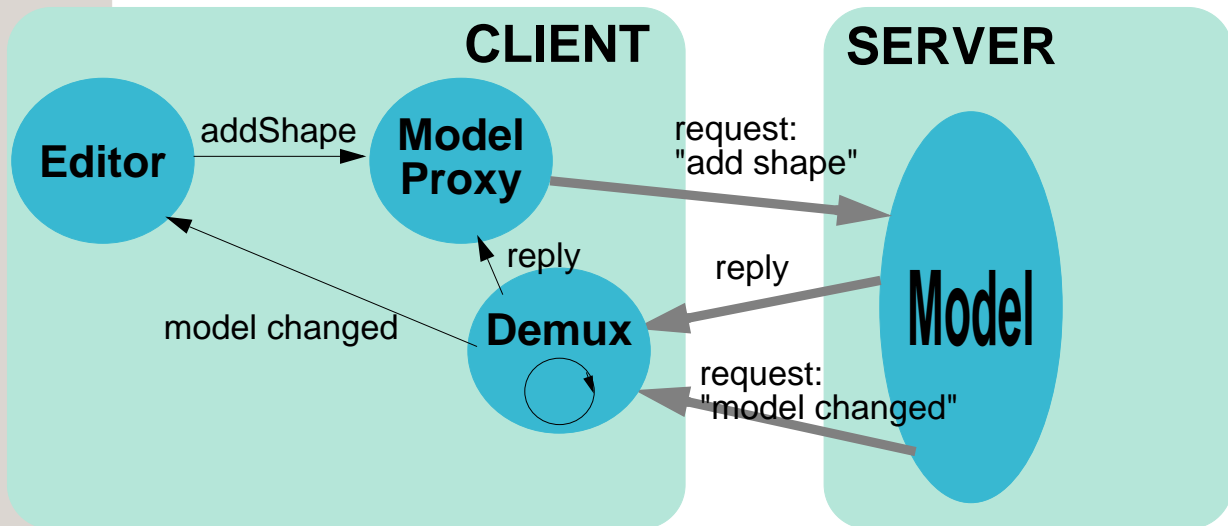
## 34 Repaint handling (local)

- "Quick hack": Model proxy send model changed requests



## 34.1 Repaint handling (Remote Observable)

- Better: "real" model sends model changed requests over the wire
- multiplex: send replies and requests on the same stream
- demultiplex the stream: separate replies from requests



## 35 Object identity in distributed programs

- when objects transferred through an ObjectOutputStream they lose their identity
- you can no longer use references to check for identity
- Solution: use an object ID
  - ◆ does not change between hosts and different runs of the program