Concurrent Systems

Nebenläufige Systeme

VII. Semaphore

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Agenda

Outline

Subject Matter

- discussion on **abstract concepts** as to unilateral and multilateral synchronisation, thus, partial and mutual exclusion
  - with the general semaphore as a measure that supports both
  - while the binary semaphore was/is intended to support the latter, only
- comprehensive differentiation of **semaphore** and **mutex**
  - in terms of the mutual exclusion aspect only, computer science folklore is right in stating disparities between the general variant and a mutex
  - but one have to be much more precise and argue with caution as far as the binary alternative is concerned

Hint (Methods v. Implementation/Entity)

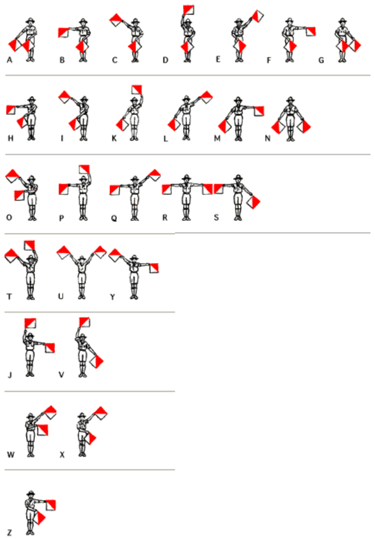
*A **binary semaphore** is a valid implementation of one of the many "mutex methods", but not that restrictive as a "mutex entity" need to be.*

- elaboration of various implementation aspects regarding both types of semaphore as well as mutex as an entity

Colloquialism

(Gr.) *sēma·pherein*, (Ger.) *Gemeinsprache*, *Redensart*

(Ger.) *Signalmast*, *Formsignal*



(Ger.) *Flaggensignal*

Outline

Concept for Cooperation and Communication

Definition (Binary Semaphore)

> *The semaphores are essentially non-negative integers; when only used to solve the mutual exclusion problem, the range of their values will even be restricted to "0" and "1". [2, p. 28]*

- jumping-off point for **sleeping lock** (Ger. *Schlafsperre*, [8, p. 9]) and, in particular, **mutex** (abbr. *mutual exclusion*)

Definition (General Semaphore)

> *It is the merit of [...] C. S. Scholten to have demonstrated a considerable field of applicability for semaphores that can also take on larger values. [2, p. 28]*

- also referred to as **counting semaphore** (Ger. *zählender Semaphor*), to control **resource allocation** in a quantitive measure

Elementary Operations

- insensitive to the distinction between binary and general semaphore is the definition of two **intrinsic primitives** [1]:

P abbr. for (Hol.) ***prolaag***; a.k.a. *down*, *wait*, or *acquire*, resp.
- decreases[1] the value of the semaphore by 1:
  - i iff the resulting value would be non-negative [2, p. 29]
  - ii non-constraining [3, p. 345]
- blocks the process iff the value is or was, resp., 0 before decrease
  - blocking processes are put on a **waitlist** associated with each semaphore

V abbr. for (Hol.) ***verhoog***; a.k.a. *up*, *signal*, or *release*, resp.
- increases[1] the value of the semaphore by 1
- as the case may be, unblocks a process blocked on the semaphore
  - which process becomes unblocked is to be regarded as unspecified

- each primitive needs to be considered as an **indivisible operation**

Hint (Waitlist)

*The **queuing discipline** rivals with planning decisions of the process scheduler and, thus, may be the cause of critical **interference**.*

[1] This does not only mean subtraction or addition, resp., in arithmetical terms.

Binary Semaphore

- **multilateral synchronisation** [5, p. 15] of interacting processes
  - the critical section is considered as a non-preemptable reusable resource that needs to be allocated indivisibly to a process to be usable correctly
  - in logical respect, the process having completed P on semaphore S is the only one being authorised to complete V on S

```
semaphore_t mutex = {1};

{
    P(&mutex);
    /* critical section */
    V(&mutex);
}
```

- default value is, normally, 1
  - block out only in the moment of a simultaneous process
  - allow full bent, else
- in case of a default value of 0
  - V must come before P

Hint (Mutex (cf. p. 14/15))

*A mutex is a **binary semaphore** that incorporates an **explicit check for authorisation** to release a critical section in the moment of V.*

General Semaphore

- **unilateral synchronisation** [5, p. 15] of interacting processes
  - used for **availability control** of entities of the following resource types:
    - i a **consumable resource** in the form of any data of any number
    - ii a **reusable resource** of limited number, e.g., a data store (buffer), any device
  - typical for, but not limited to, producer/consumer systems
- also as noted previously [5, p. 15], this art of synchronisation means:
  - logical
    - coordination as indicated by a particular "role playing"
    - e.g., in order to proceed, a "data consumer" depends on the data to be made available by a "data producer"
  - conditional
    - coordination as indicated by a condition for making progress
    - e.g., in order to proceed, a "data producer" depends on the store available for data handling
    - in the end, the data store will have to be deallocated and, thus, made available again by the "data consumer"
- from this it follows that P and V applied to the same semaphore S must have to be accomplishable by different processes, normally
  - which makes the big difference to a binary semaphore or mutex, resp.

Consumable Resource


```
semaphore_t data = {0};

void producer() {
    for (;;) {
        /* data released */
        V(&data);
    }
}

void consumer() {
    for (;;) {
        P(&data);
        /* data acquired */
    }
}
```

- default value is 0
  - P must block out only if there is no data
  - V indicates more data
- calling sequence
  - V must be actable independent of P
  - in order to complete, P depends on V
  - ↪ beware of an **overflow** of the values margin
- usually, producer and consumer are <u>different</u> interacting processes
  - in case of one and the same process, the number of a completed V must exceed the number of a completed P in order to prevent deadlock
  - $\#V > \#P$, which implies a path $V \rightarrow P$ (i.e., V "happens before" P)

Reusable Resource


```
semaphore_t store = {N};

void producer() {
    for (;;) {
        P(&store);
        /* store acquired */
    }
}

void consumer() {
    for (;;) {
        /* store released */
        V(&store);
    }
}
```

- default value is $N \geq 0$
  - P must block out only if there is no store
  - V indicates more store
- calling sequence
  - V must be actable independent of P
  - in order to complete, P depends on V
  - ↪ beware of an **overflow** of the values margin
- as to interacting processes in the line of producer and consumer, the same applies as mentioned before: $\#V > \#P$
- in other cases: $\#V \leq \#P$, must be completed by the <u>same</u> process

Availability Control in Practice — Bounded Buffer

Hint (Bounded Buffer)

A means of managing an unlimited number of consumable resources on the basis of a limited number of reusable resources.

```
semaphore_t data = {0}, store = {N};     /* N > 0 */

void producer() {
  for (;;) {
    P(&store);
    /* store acquired */
    /* data released */
    V(&data);
  }
}
```

```
void consumer() {
  for (;;) {
    P(&data);
    /* data acquired */
    /* store released */
    V(&store);
  }
}
```

- indisputable classic in cooperation and communication of processes
  - simply a merge of the semaphore use pattern discussed as before
  - **transverse application** of P and V to a pair of general semaphores

Semaphore v. Mutex I — Technical Level

Hint

*Checking **authorisation** for release of a critical section in that very moment is improper for a general semaphore, optional for a binary semaphore, and may be demanded for a mutex (cf. p. 15).*

demanded
- a **mutex entity** ensures that the release of critical section CS will succeed only for the process having acquired CS
- by extending a binary semaphore, P will have to record and V will have to check ownership of CS

improper
- P and V on a **general semaphore** must be accomplishable in particular also by different processes
- this is prevented by a mutex entity—but not by a mutex

optional
- basically, a **binary semaphore** may be implemented by a general semaphore S, with $S \leq 1 \Rightarrow$ never a mutex entity
- values $S > 1$ must be prevented either by the use pattern or by the implementation of P and V

- if **authorisation fails**, the process attempting to release CS should be aborted—in kernel mode, the computing system must be halted...

Semaphore v. Mutex II — Conceptual Level

Hint (Computer Science Folklore)

A semaphore can be released by any process.

- incomplete or rough, if not broad-brush, phrase that must be regarded with suspicion—one have to distinguish between semaphore types[2]
  - strictly, essence of this phrase is **requirement** for a general semaphore
  - strictly as well, it is merely an **option** for a binary semaphore
    - in logical respect, a binary semaphore may not be released by any process
    - in physical respect, this however is not a must for any implementation

Hint (Computer Science Folklore)

A mutex can be released only by the process having it acquired.

- a phrase that is slanted towards only one aspect as to the leastwise twofold non-uniform common understanding about a mutex:
  - i a category of **methods** for ensuring mutual exclusion *or*
  - ii the **implementation** of one of these methods in terms of an **entity**[2]

[2] see also p. 38

Hierarchic Placement — "Uses" [4] Process Instances

- the **standby position** of a process within P is passive, normally
  - "blocks the" or "unblocks a", resp. (cf. p. 8), process means rescheduling
  - if so, both may also entail context switching—"may" because:
    - P – if no further process is ready to run, the ***idle loop*** becomes active
      - in that case, the blocking process likewise may fade to the ***idle process***
      - thus, doing without a dedicated ***idle-process instance*** and context switch
    - V – if there is a waiting process, it will be set "ready to run" (cf. [9, p. 28])
      - in that case, ***priority violation***[3] must be prevented (scheduling discipline!)
      - thus, the current process may defer to a prior-ranking one: context switch
  - all this makes P and V programs of the operating system machine level
- P and V relies on **process management** of the operating system
  - one have to put the current process asleep and get a sleeping process up
  - in functional terms, however, P and V need not be system calls
  - in non-functional terms, P and V should be close to the **scheduler**
    - by settling P and V in the address space of the operating-system kernel *or*
    - by making scheduler functions available through "strawweight" system calls

[3] If at least one of the processes on the waitlist is of higher-priority than the current process but will not become "ready to run" or allocated the processor.

Waitlist Association

- in order to aid V, processes blocked by P at a semaphore are entered on a waitlist in either logical or physical means
  - logical
    - to block, a **blocked-on mark** is stored in the process descriptor
    - to unblock, a process-table walk looks for that mark
    - ↪ constant (P) and variable but bounded above (V) run-time
    - ↪ blocked-on mark is a "magic" address, no extra attributes
  - physical
    - to block, the process descriptor joins a **queue data structure**
    - to unblock, a process descriptor is removed from that structure
    - ↪ variable but bounded above (P) and constant (V) run-time
    - ↪ additional queue attribute of the semaphore data structure
- desirable is to have the waitlist queuing discipline in compliance with the process scheduling discipline: **freedom of interference**
  - a characteristic by means of which **priority violation** will be prevented
  - usually, this excludes straightforward queuing disciplines such as FCFS

Hint (Process-Table Walk—Conformance to Scheduling)

Part of the scheduler, lookup function to locate a process descriptor on the basis of the blocked-on mark as search key.

Indivisible Operation I

Rationale

- in the absence of simultaneous processes, the implementation of a semaphore could be as simple as follows:

```
1 void prolaag(semaphore_t *sema) {
2     if (!claim(sema))    /* at the moment, unavailable */
3         sleep(&sema->wand);
4 }
5
6 void verhoog(semaphore_t *sema) {
7     if (unban(sema))    /* as from now, available */
8         rouse(&sema->wand);
9 }
```

  - whereat `claim` decreases and `unban` increases the value of the semaphore according to binary or general, resp., characteristic[4]
- but, assuming that the presence of simultaneous processes is possible, this implementation shows a **race condition** ⇝ ***lost wakeup***
  - 3: while going to sleep, i.e. being "sleepy", the process gets delayed
  - 7–8: but in good faith of a sleeper, the "sleepy" process may be missed

[4] The implementation of these helper functions will be revealed later.

Indivisible Operation II

Scope

- P and V itself constitute a **critical section**, likewise, that must be protected in order to function correctly
  - protection should be constructed **per semaphore instance**, not P/V

```
void prolaag(semaphore_t *sema) {
    atomic *sema = {
        if (!claim(sema))
            sleep(&sema->wand);
    }
}

void verhoog(semaphore_t *sema) {
    atomic *sema = {
        if (unban(sema))
            rouse(&sema->wand);
    }
}
```

Deadlock Prevention

Provided that protection of the critical section on the P side is not deregulated, the V side will never complete and, thus, will never cause unblocking of a process:

- the right location for deregulation is `sleep`
- after the process was marked sleeping

- as a process will have to block inside a critical section, **deregulation of protection** is indispensable for the period the process is blocked

Shallows

Implementation

- protection of the P/V pair against simultaneous processes sharing a semaphore follows either the blocking or non-blocking paradigm
  - blocking
    - inhibit FLIH[5], postpone SLIH[5], or lock process
    - problem-specific construction of an ***enter***/***leave*** pair
    - ↪ coming right up next in this lecture (cf. p. 22*ff.*)
  - non-blocking
    - fall back on the elementary operations of the ISA level
    - problem-specific construction of P and V
    - ↪ coming up as a case study in the context of LEC 10/11
- more detailed analysis of the "atomic" version of P reveals another problem: **overtaking** of an aroused process
  - upon return from *sleep* a formerly blocked process may complete P by mistake, joining a process in the critical section to be protected by P
  - this is because completion of V also opens the door for any process, not only for a process having been blocked at the semaphore
  - ↪ if applicable, aroused processes have to **retry claiming**: **`if`** ↦ **`while`**
- not least, concurrency had to be constricted to no more than what is absolutely necessary: reflect on `claim`/`sleep` and `unban`/`rouse`

[5] abbr. for *first-* or *second-level interrupt handling*, resp.

Outline

Semaphore Data Type

```
typedef volatile struct semaphore {
    int gate;              /* value: binary or general */
    wand_t wand;           /* protective shield */
} semaphore_t;
```

- purpose of "wand" (Ger. *Zauberstab*) is to **safeguard** the semaphore operations in various respect
  - i protect P and V against simultaneous processes
  - ii give leeway for protection variants (cf. p. 20)
- a wand that takes care of **mutual exclusion** techniques by means of locks [8], for example, could be the following:

```
typedef volatile struct wand {
    lock_t clue;           /* protects P or V, resp. */
    event_t wait;          /* list of sleeping processes */
} wand_t;
```

↪ becoming acquainted with other wands is content of future lectures...

P and V Safeguarded

Mutual Exclusion

```
1  void prolaag(semaphore_t *sema) {
2      enter(&sema->wand);      /* avert overlapped P or V */
3      lodge(sema);             /* raise claim to proceed */
4      when (!avail(sema))      /* check for process delay */
5          sleep(&sema->wand);  /* await wakeup signal */
6      leave(&sema->wand);      /* allow P or V */
7  }
8
9  void verhoog(semaphore_t *sema) {
10     enter(&sema->wand);      /* avert overlapped P or V */
11     if (unban(sema))         /* release semaphore */
12         rouse(&sema->wand);  /* cause wakeup signal */
13     else                     /* no sleeping process... */
14         leave(&sema->wand);  /* allow P or V */
15 }
```

- exercise caution in the analysis of these program statements:
  - 4 ■ if applicable, "when" takes care of overtaking processes
  - 11–12 ■ if applicable, search for sleepers happens unconditionally
    - in case of (i) logical waitlist and (ii) strict binary semaphore

Acquire and Release Semaphore I

- load/store-based implementation for a **binary semaphore**:

```
inline int lodge(semaphore_t *sema) {
    return 42;
}

inline bool avail(semaphore_t *sema) {
    return (sema->gate == 0) ? false : !(sema->gate = 0);
}

inline bool unban(semaphore_t *sema) {
    return (sema->gate = 1) && exist(&sema->wand);
}
```

- note that the semaphore value alone shows no indication of processes that potentially await a reveille (Ger. *Wecksignal*) as to this very semaphore
- only an explicit waitlist scan sheds light on that ⇝ `exist`

- also note the persisting sensitivity to simultaneous processes: `avail`
  - use within a safeguarded program section is assumed...

Acquire and Release Semaphore II

- enumerator-based implementation for a **general semaphore**:

```
inline int lodge(semaphore_t *sema) {
    return sema->gate--;
}

inline bool avail(semaphore_t *sema) {
    return sema->gate >= 0;
}

inline bool unban(semaphore_t *sema) {
    return (sema->gate++ < 0);
}
```

  - note that the absolute value of a "negative semaphore" gives the number of processes awaiting a reveille as to this very semaphore
  - thus, there is no need for an explicit waitlist scan ☺

- also note the persisting sensitivity to simultaneous processes: --/++
  - use within a safeguarded program section is assumed...

Overtaking Lane: Semaphore Handicap (cf. p. 20)

- in contrast to the shown general semaphore, a roused process has to **recheck his waiting condition** in *P*

- **general semaphore**
  - **#define when if**
    - susceptible to erroneous rouse
  - overtaking impossible
    - $gate \leq 0$ when a process aroused
    - rival process in *P* causes $gate < 0$
  - ↪ will be forced to sleep
  - ↪ aroused process may proceed

- **binary semaphore**
  - **#define when while**
    - unsusceptible to erroneous rouse
  - overtaking possible
    - $gate = 1$ when a process aroused
    - rival process in *P* causes $gate = 0$
  - ↪ is allowed to continue ☹
  - ↪ aroused process has to wait

Hint (erroneous rouse)

Caused by misuse of V or by forced and uncontrolled unblocking of a process that went to sleep in P. Both are programming errors: the former at (semaphore) application level, the latter at system level.

Special Process Management — Prevent Lost Wakeup

```
1  inline void sleep(wand_t *wand) {
2      catch(&wand->wait); /* disclose process to V */
3      leave(wand);        /* allow P or V */
4      coast();            /* take a break */
5      enter(wand);        /* apply for return to P */
6  }
7
8  inline void rouse(wand_t *wand) {
9      leave(wand);        /* allow P or V */
10     cause(&wand->wait); /* signal end of break */
11 }
```

- constrict concurrency to no more than what is absolutely necessary:
  - 2 ▪ endorse interest of the current process of upcoming dormancy
  - 3 ▪ soon dormant process was made known, deregulate *P* safeguard ←
  - 4 ▪ transition to dormant state: rescheduling, context switch or idleness
  - 5 ▪ apply for return to safeguarded *P*
  - 9 ▪ dormant processes could be available, deregulate *V* safeguard ←
  - 10 ▪ annulment of dormant state: rescheduling, context switch

General Process Management — Event Handling

catch
- has two variants, depending on the waitlist model (cf. p 17):
  - i store of a blocked-on mark in the process descriptor *or*
  - ii enqueue of the process descriptor into a queue data structure
- variant (i) writes to an own data structure of the current process, while variant (ii) manipulates a shared data structure
- signalises upcoming blocking (dormancy) of registered process

coast
- blocks the current process, reschedules the processor, and either performs a context switch or runs through the idle loop
  - manipulates a shared data structure (ready list)
  - performs the queuing function of the queue-based *catch*
- eventually returns when the blocking condition was nullified

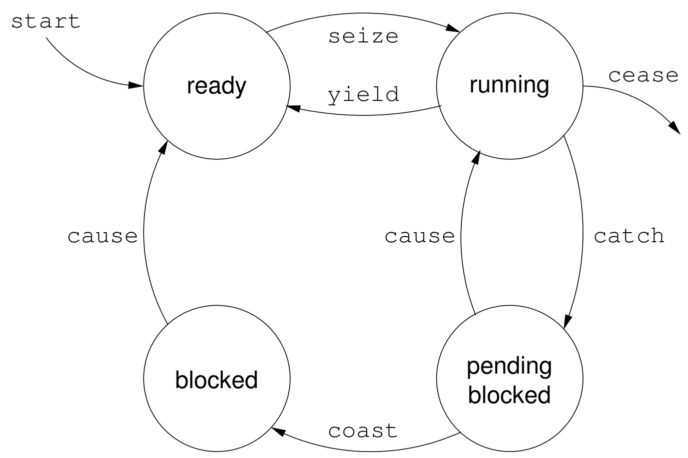
cause
- unblocks the next registered process, if any, found by means of a (i) process-table walk or (ii) dequeue operation
  - manipulates a shared data structure (ready list)
- if need be, the current process defers to a prior-ranking process

Hint (Idle State (cf. p. 16 and p. 39))

The last process blocked may find itself on the ready list. Same may happen to the "sleepy process" as coast runs deregulated to P/V.

Process States and State Transitions

cf. [9, p. 27]



- ready ↔ running ■ scheduler
- blocked → ready ■ iff *effective signalling* (V), i.e., waiting process
- running ↔ pending ■ doze (P →), *effective signalling* (← V)
- pending → blocked ■ deep sleep (P), no overlapping V

Semaphore Gatekeeper

(Ger.) *Schrankenwärter, Türhüter*

- as there is no single solution to protect P and V adequately, the wand attribute symbolises intention to application orientation
  - depending on the mode of operation or use case, the wand acts differently
  - assuming that processing elements are multiplexed [7, p. 5], then:

```
inline void enter(wand_t *wand) {
    avert(IRQ); lock(&wand->clue);
}

inline void leave(wand_t *wand) {
    unlock(&wand->clue); admit(IRQ);
}
```

Kernel-level Safeguard:

1. disable IRQ
2. lock critical section

- wand capability depends on the "type of exclusion" in relation to the required characteristics of the operating system machine level:
  - partial ■ processor **multiplexing** ⇝ interrupt (IRQ) control
  - mutual ■ processor **multiplication** ⇝ process lock
- combination of both is optional, not mandatory, and problem-specific
  - depends on the degree of parallelism (a) allowed for by the application use case and (b) made possible by the ISA level

Specialisation of a Binary Semaphore

Mutex (cf. p. 14/15)

- given the concept of a binary semaphore, implementation of a **mutex** is straightforward and, absolutely, no black magic
  - a mutex data structure is composed of two parts:
    - i a binary semaphore used to actually protect the critical section *and*
    - ii a handle that uniquely identifies the process having acquired the mutex[6]
  - given such a structure, let the following two functions be defined:
    - acquire – performs the P and registers the current process as owner
    - release – conditionally unregisters the owner and performs the V
      – in case of a wrong owner, the current process or kernel panics
- a corresponding **data type** may be laid out as follows:

```
typedef volatile struct mutex {
    semaphore_t sema;   /* binary semaphore */
    process_t *link;    /* owning process or 0 */
} mutex_t;
```

[6] At kernel level, the handle is the pointer to the process descriptor of the process instance. At user level, it is the process identification.

Acquire and Release Mutex

```
extern void panic(char*) __attribute__ ((noreturn));

void acquire(mutex_t *mutex) {
    P(&mutex->sema);          /* lockout */
    mutex->link = being(ONESELF);    /* register owner */
}

void release(mutex_t *mutex) {
    if (mutex->link != being(ONESELF)) /* it's not me! */
        panic("unauthorised release of mutex");

    mutex->link = 0;          /* deregister owner */
    V(&mutex->sema);          /* unblock */
}
```

- release of a mutex by an **unauthorised process** is a serious matter
  - presumably, the non-sequential program contains a **software fault** (bug)
  - returning an error code is no option, as one cannot rely on error checking
  - any other than "raising a non-maskable exception" is a botch job...

Outline

Résumé

- fundamental concept for cooperation and communication
  - binary and general/counting semaphore, intrinsic primitives P and V
  - correlation to unilateral and multilateral synchronisation
  - differentiation as to mutex (methods v. implementation/entity):

Hint

*A **binary semaphore** is a valid implementation of one of the many "mutex methods", but not that restrictive as a "mutex entity" need to be.*

  - hierarchic placement at operating system machine level
- characteristics important in functional and non-functional terms
  - logical or physical waitlist, conformance to the scheduling discipline
  - deregulation of the protection of P against simultaneous processes
  - further shallows such as overtaking of unblocked processes in P:

Hint

Constrict concurrency to no more than what is absolutely necessary.

- not least, basic approaches and sketches of an implementation. . .

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Extent of Critical Section

Binary Semaphore devoid of Waitlist[7]

```
1        ...
2        movl   16(%esp), %edi
3        leal   4(%edi), %esi
4        jmp    LBB0_2
5   LBB0_1:
6        movl   _life, %eax
7        movl   %esi, 4(%eax)
8        movl   %esi, (%esp)
9        calll  _unlock
10       calll  _coast
11  LBB0_2:
12       movl   %esi, (%esp)
13       calll  _lock
14       cmpl   $0, (%edi)
15       je     LBB0_1
16       movl   $1, (%edi)
17       movl   %esi, (%esp)
18       calll  _unlock
19       ...
```

- let the sequence of instructions within P be as follows:
  - 2 point at semaphore
  - 3 point at lock structure
    - address is blocked-on mark
  - 12–13 apply for P protection
  - 14–15 check binary semaphore S_b
  - 16–18 unoccupied, take S_b
    - quit P protection, done
  - 5 occupied, S_b already taken
  - 6 point at process structure
  - 7 define blocked-on mark
  - 8–9 deregulate P protection
  - 10 fall asleep, dream about V
- locking overhead when unoccupied
  - net worth of about 5 instructions ☹
  - ↪ non-blocking synchronisation ☺

[7] Take a sledgehammer to crack a nut...

Semaphore v. Mutex III

cf. p. 14/15

Commonalities and differences as to their possible **internal states**.

- general semaphore S_g:
  - positive: $N > 0$ processes will complete $P(S_g)$ without blocking
  - zero: $P(S_g)$ will block the running process on the waitlist of S_g
  - negative: $P(S_g)$ will block the running process on the waitlist of S_g
    - $|N|$ processes are blocked on the waitlist of S_g
- binary semaphore S_b:
  - not taken: exactly one process will complete $P(S_b)$ without blocking
    - the very process becomes **logical owner** of S_b
  - taken: $P(S_b)$ will block the running process on the waitlist of S_b
    - $V(S_b)$ should be performed only by the logical owner of S_b
- mutex entity M: let A be *acquire* and let R be *release*
  - not owned: exactly one process will complete $A(M)$ without blocking
    - the very process becomes **physical owner** of M
  - owned: $A(M)$ will block the running process on the waitlist of M
    - $R(M)$ can succeed only for the physical owner of M

Idle State

cf. p. 28

- principle pattern of a scheduler function to block a process
  - called by coast (cf. p. 27) and other functions to pause computation

```
1   void block() {
2       process_t *next, *self = being(ONESELF);
3
4       while (!(next = elect(hoard(READY))))
5           relax();                 /* no ready to run... */
6
7       if (next != self) {          /* must relinquish */
8           self->state = BLOCKED;   /* vacate processor */
9           seize(next);             /* resume elected */
10      }
11      self->state = RUNNING;       /* occupy processor */
12  }
```

- 4 choose next process to be dispatched to the processor
- 5 ready list is empty, so the running process fades to the idle process
- 7 as the case may be, the running process may be allowed to continue:
  - i the idle/running process found itself ready-to-run on the ready list *or*
  - ii the running process, sent to sleep due to P, was roused due to V (p. 27)